



Panasas Storage System – Integration into a Cluster

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Agenda



- HPC requirements for storage
- Shared storage – promise and challenges
- Need for new storage architecture, object-based storage
- Panasas implementation of object storage architecture
- Panasas integration into a cluster with existing storage
- Panasas integration at Stanford University – Steve Jones, Stanford
- Demo – Alex Krimkevich, Panasas
- Questions and answers

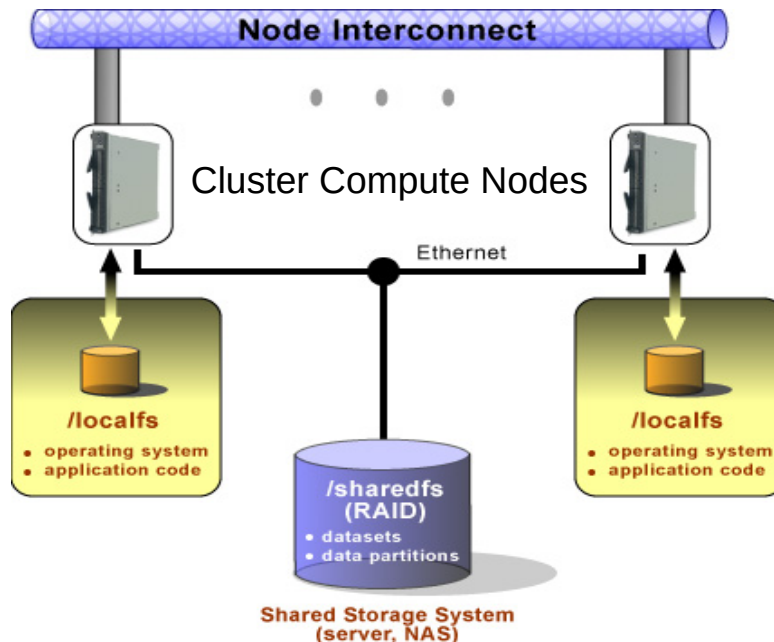
Storage Requirements for HPC, Rocks



- Performance
 - High read concurrency for parallel application and data sets
 - High write bandwidth for memory checkpointing, interim and final output
- Scalability
 - More difficult problems typically means larger data sets
 - Scaling cluster nodes requires scalable IO performance
- Management
 - Single system image maximizes utility for user community
 - Minimize operations and capital costs (80% of storage TCO)

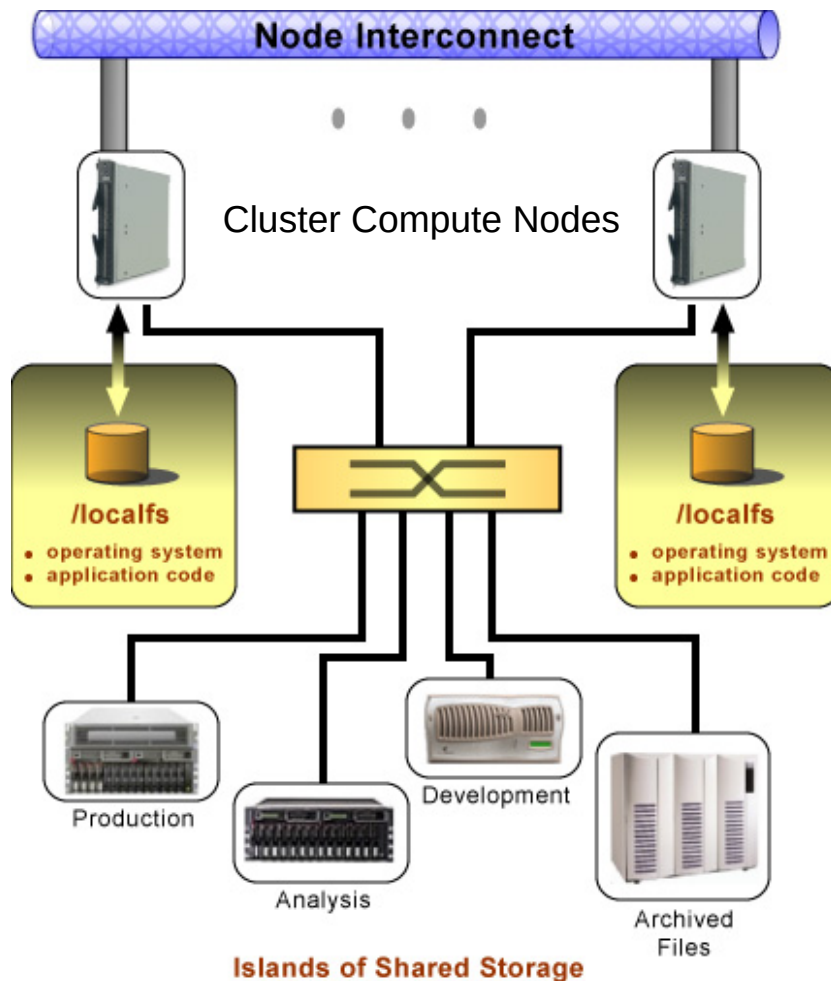
High demand on storage systems

Shared Storage: The Promise



- Shared storage cluster computing
 - Compute anywhere model
 - Partitions available globally; no replicas required (shared datasets)
 - No data staging required
 - No distributed data consistency issues
 - Reliable checkpoints; application reconfiguration
 - Results gateway
 - Enhanced reliability via RAID
 - Enhanced manageability
 - Policy-based management

Shared Storage Challenges



- Performance, scalability & management
 - Single file system performance limited
 - Multiple volumes and mount points
 - Manual capacity and load balancing
 - Large quantum upgrade costs

Motivation for New Architecture

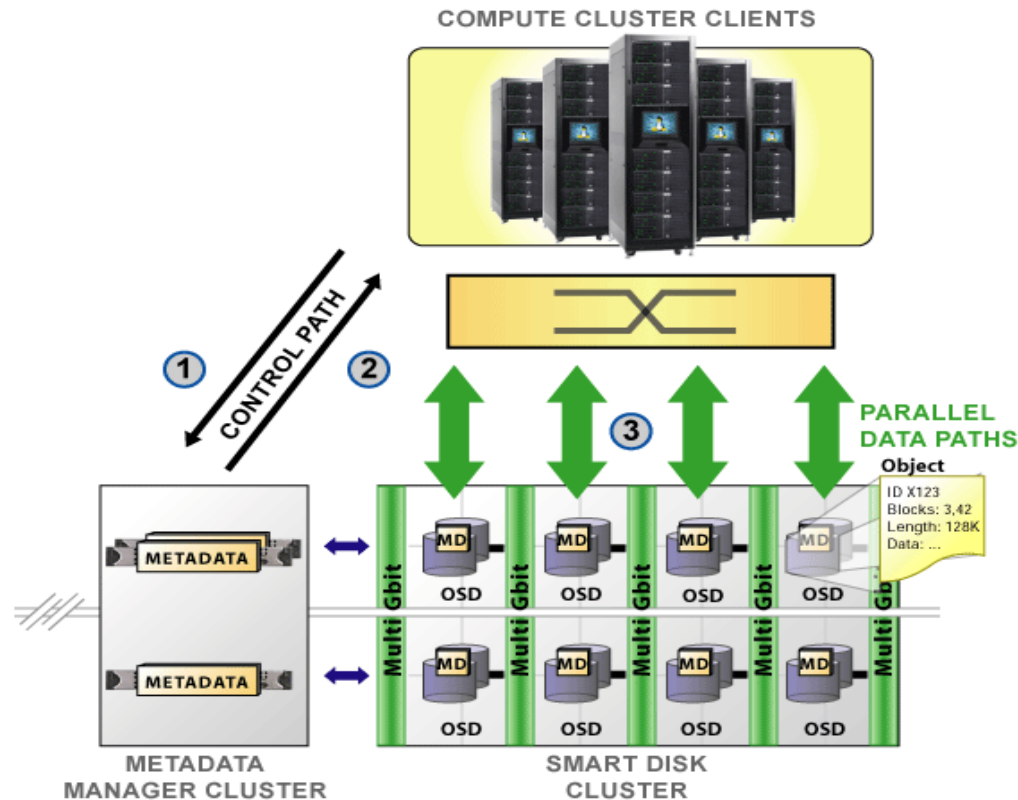


- A highly scalable, interoperable, shared storage system
 - Improved storage management
 - Self-management, policy-driven storage (i.e. backup and recovery)
 - Improved storage performance
 - Quality of service, differentiated services
 - Improved scalability
 - Of performance and metadata (i.e. free block allocation)
 - Improved device and data sharing
 - Shared devices and data across OS platforms

Next Generation Storage Cluster



OBJECT STORAGE ARCHITECTURE



HOW IT WORKS

Object-based Storage is an emerging networked storage standard that delivers unprecedented performance and scalability within a single management domain.

- 1 Client processes make independent file requests (e.g., stat, open) of the clustered metadata manager
- 2 Authenticated clients receive object maps and capabilities
- 3 Clients use their credentials to perform parallel direct-to-disk requests across the GbE switching network

Panasas ActiveScale File System

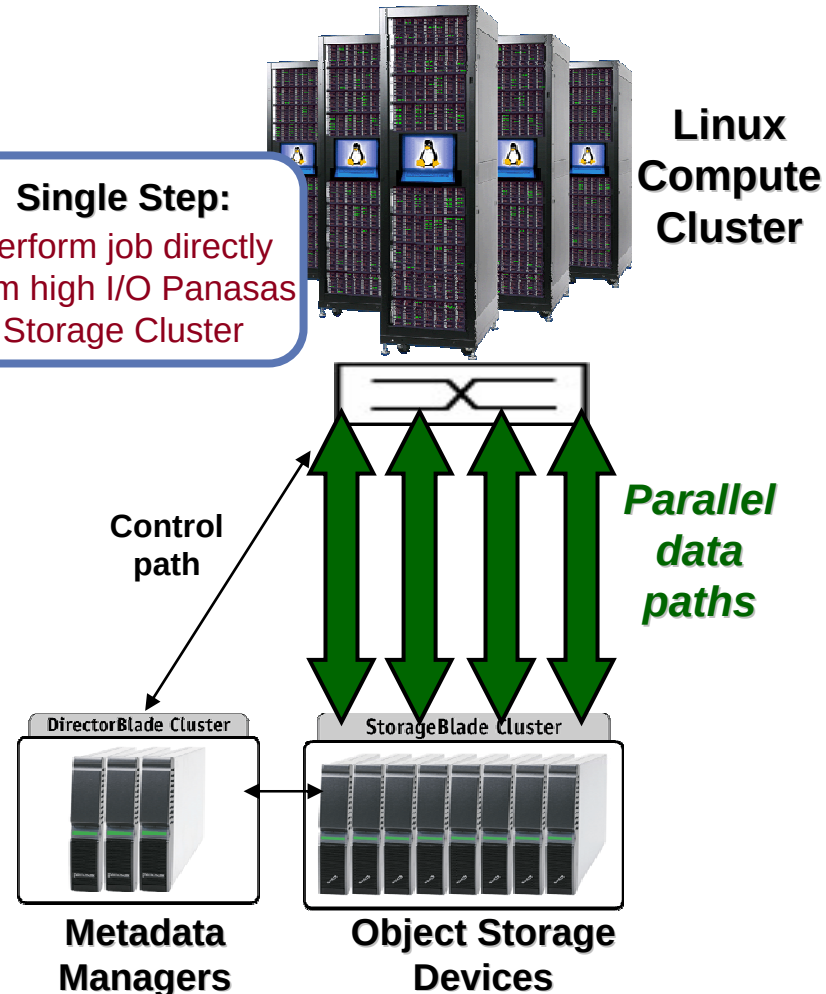
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BY



panasas

- Scalable performance
 - Offloaded data path enables nodes to access disks directly
 - Scale nodes, network, and capacity
 - As capacity grows, performance grows
- Simplified and dynamic management
 - Robust, shared file access
 - Seamless growth within single namespace
- Integrated HW/SW solution
 - Optimizes performance and manageability
 - Ease of integration and support

Single Step:
Perform job directly
from high I/O Panasas
Storage Cluster



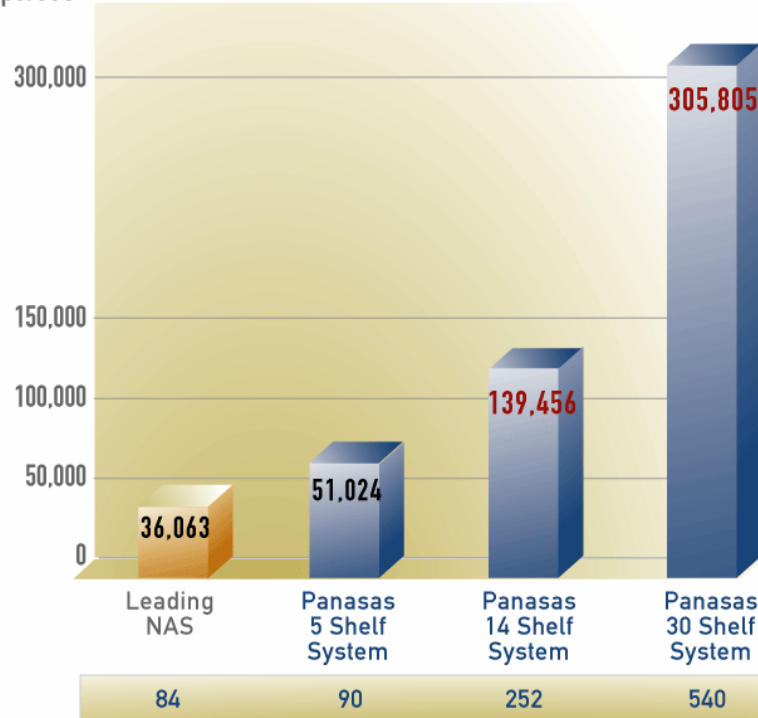
Panasas Storage Cluster

Industry-Leading Performance



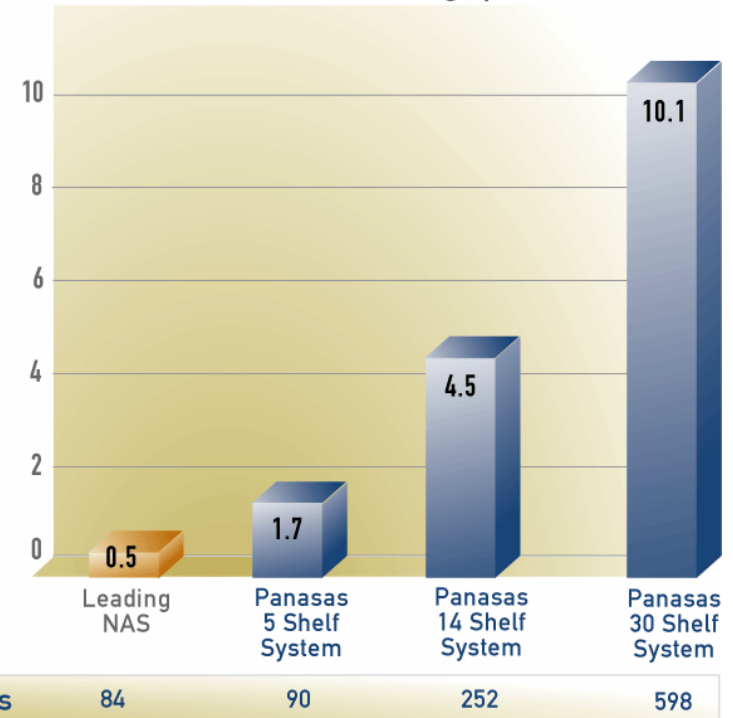
SPECsfs97_R1
Ops/sec

Random I/O



GB/sec

Data Throughput



Breakthrough data throughput AND random I/O

Bottlenecks in the IT Infrastructure

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ROCKS



CLIENTS (End Users)



Clients

- Clients waiting to access job results
- Scalability problems as users increase
- Hindrance to collaboration and sharing data

Cluster

- Cluster waiting for I/O complete
- Reduced cluster utilization
- Scalability problems as nodes increase

BOTTLENECK

Backup / Restore

- Backup interfering with other processes
- Backup window pressures
- Inability to do fast backups and restores

TRADITIONAL
NETWORK STORAGE

BOTTLENECK

BOTTLENECK



BACKUP / RESTORE



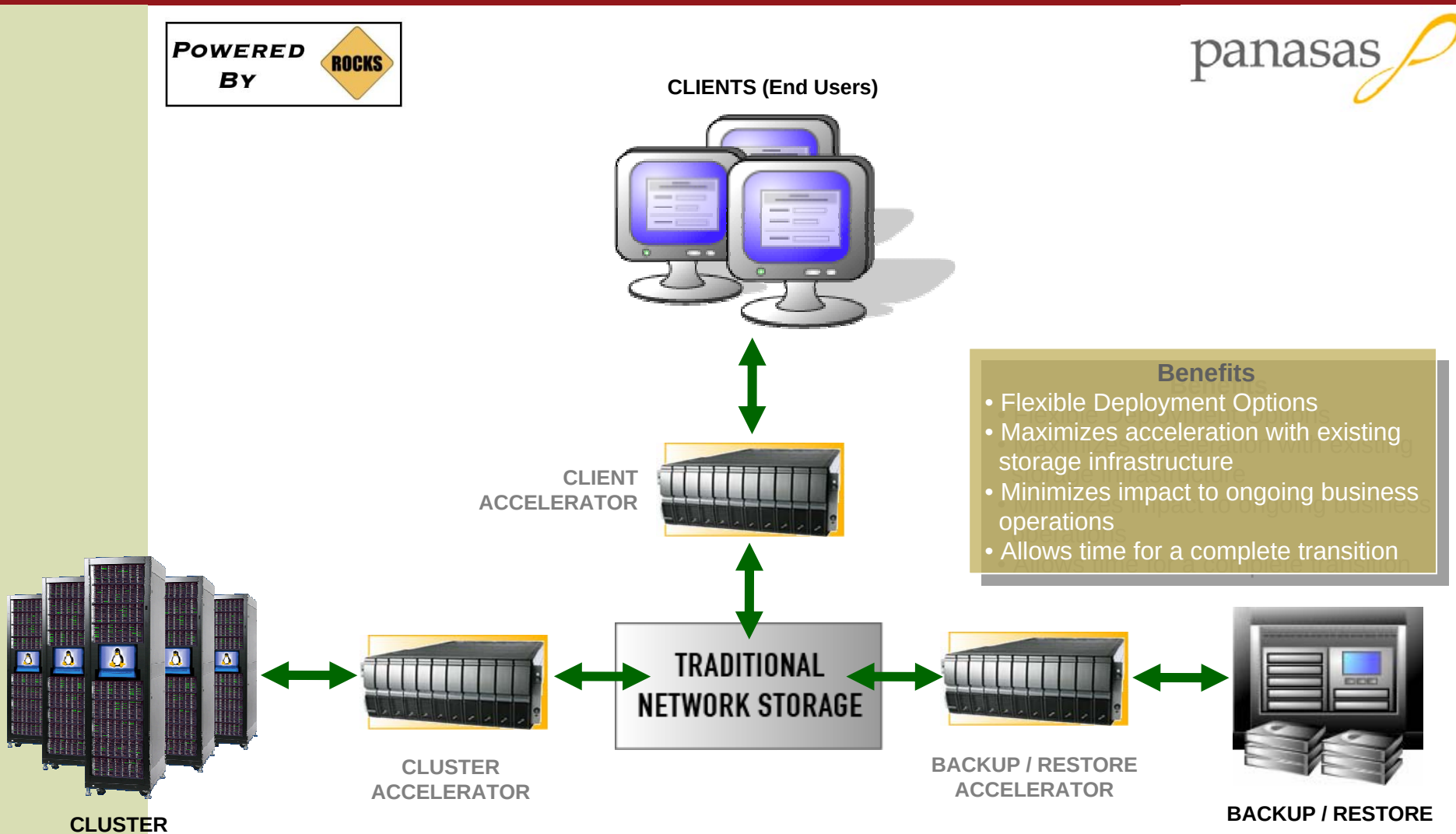
CLUSTER

Accelerators To Boost Performance



- **Cluster Accelerator**
 - Parallel reads, writes
 - Higher cluster utilization, more jobs processed
 - Leverages existing storage assets
- **Client Accelerator**
 - Fast access to job results for end users
 - Promotes efficient collaboration and independent analysis
- **Backup / Restore Accelerator**
 - Staging area for fast backups / restores
 - Offloads backup from production storage
 - Works with backup / restore applications

Maximum Acceleration with Existing Storage



Complete Acceleration

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ROCKS

panasas

CLIENTS (End Users)



COMPLETE
IMPLEMENTATION OF
PANASAS STORAGE



BACKUP / RESTORE



CLUSTER

Benefits

- Removes bottlenecks of traditional storage
- Minimizes "data motion"
- Compute from home directories
- Output accessible from everywhere
- Data processing pipeline
- Output from one process is input to next process

Summary



- HPC puts unprecedented demands on storage systems
- Object storage architecture fulfills the promise of shared storage
 - Manageability (80% of storage TCO)
 - Scales performance in concert with capacity
 - Implemented by Panasas with industry leading performance
- Integration of Panasas storage with clusters is non-disruptive
 - Deploy with existing storage infrastructure as Accelerators
 - Cluster, Client, Backup / Restore
 - Gain immediate benefits
 - Transition to complete deployment over time
- Key to productivity with HPC – eliminate the bottlenecks to storage systems, simplify management



Panasas Storage System – Integration into a Cluster at **Stanford University**

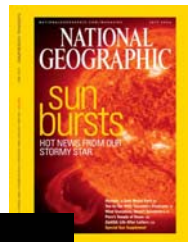
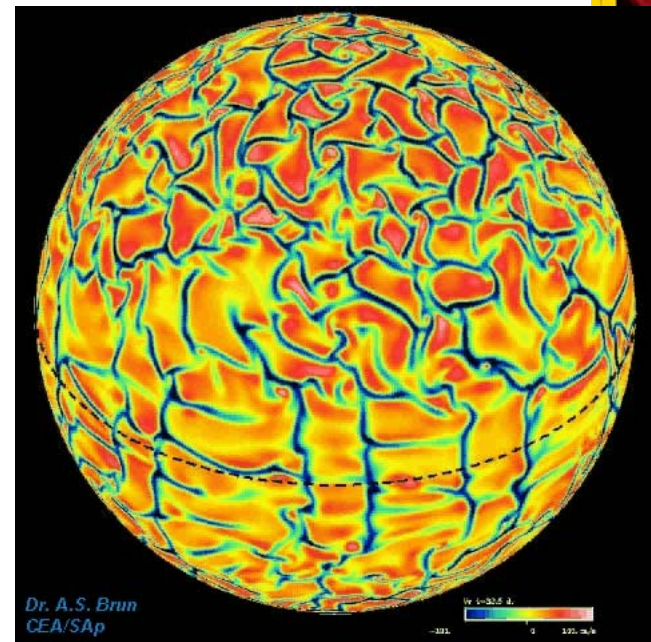
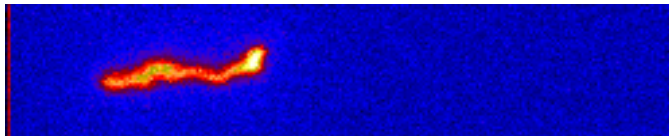
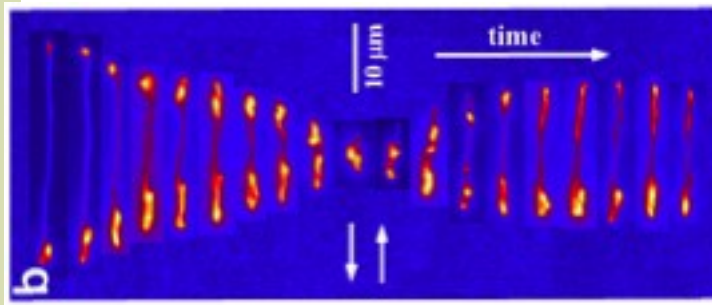
Steve Jones

Technology Operations Manager
Institute for Computational and Mathematical Engineering
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The Research



MOLECULES TO PLANETS !



Research Groups



- Flow Physics and Computation
- Aeronautics and Astronautics
- Chemical Engineering
- Center for Turbulence Research
- Center for Integrated Turbulence Simulations
- Thermo Sciences Division

Funding

- Sponsored Research
(AFOSR/ONR/DARPA/DURIP/ASC)

Active Collaborations with the Labs



Buoyancy driven instabilities/mixing - CDP for modeling plumes
(Stanford/SNL)

LES Technology - Complex Vehicle Aerodynamics using CDP
(Stanford/LLNL)

Tsunami modeling - CDP for Canary Islands Tsunami Scenarios
(Stanford/LANL)

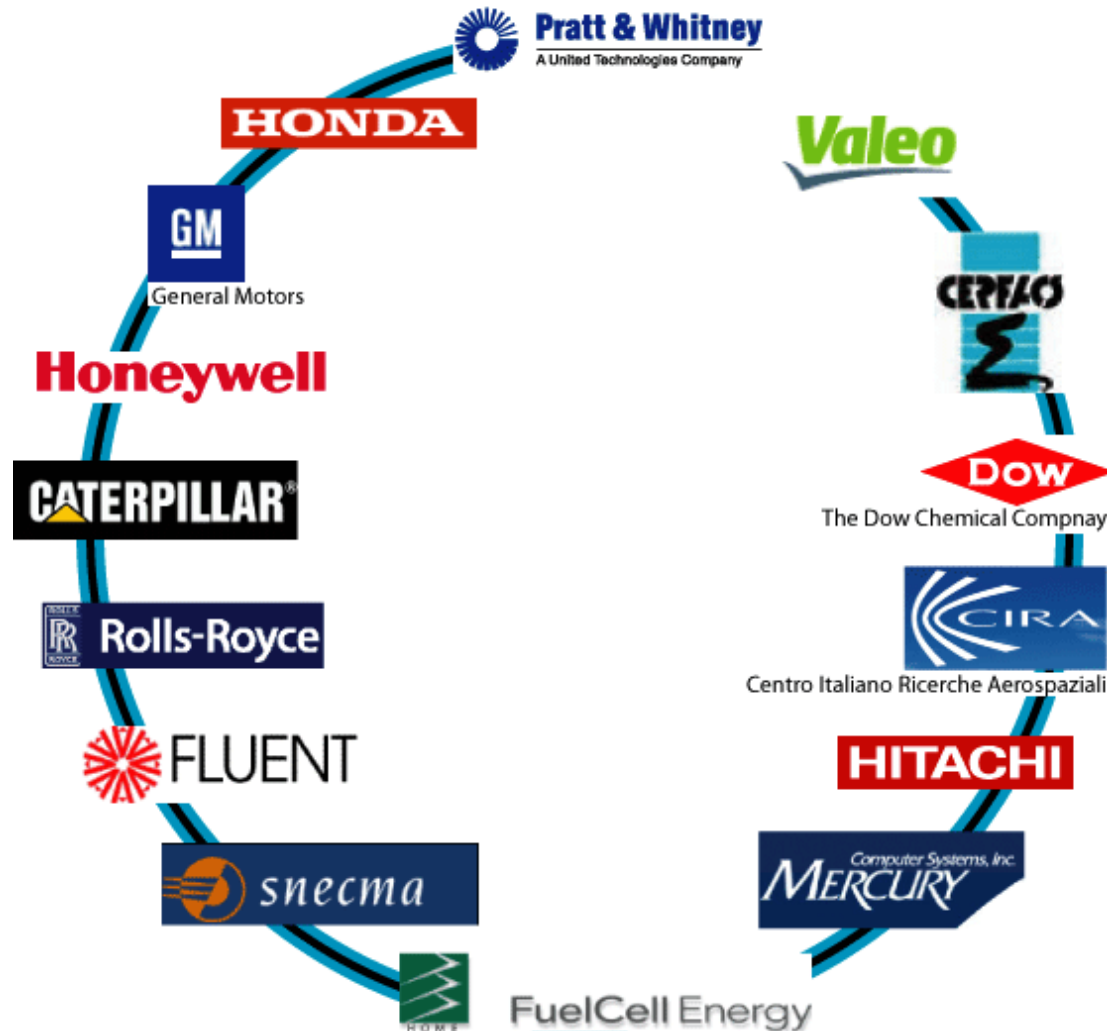
Parallel I/O & Large-Scale Data Visualization - UDM integrated in CDP
(Stanford/LANL)

Parallel Global Solvers - HyPre Library integrated in CDP
(Stanford/LLNL)

Parallel Grid Generation - Cubit and related libraries
(Stanford/SNL)

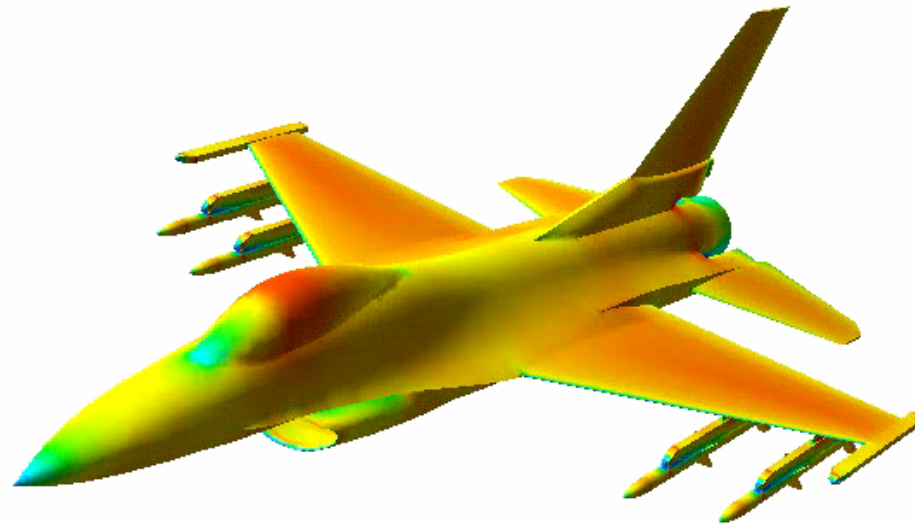
Merrimac - Streaming Supercomputer Prototype
(Stanford/LLNL/LBNL/NASA)

Affiliates Program



F-16 Simulation

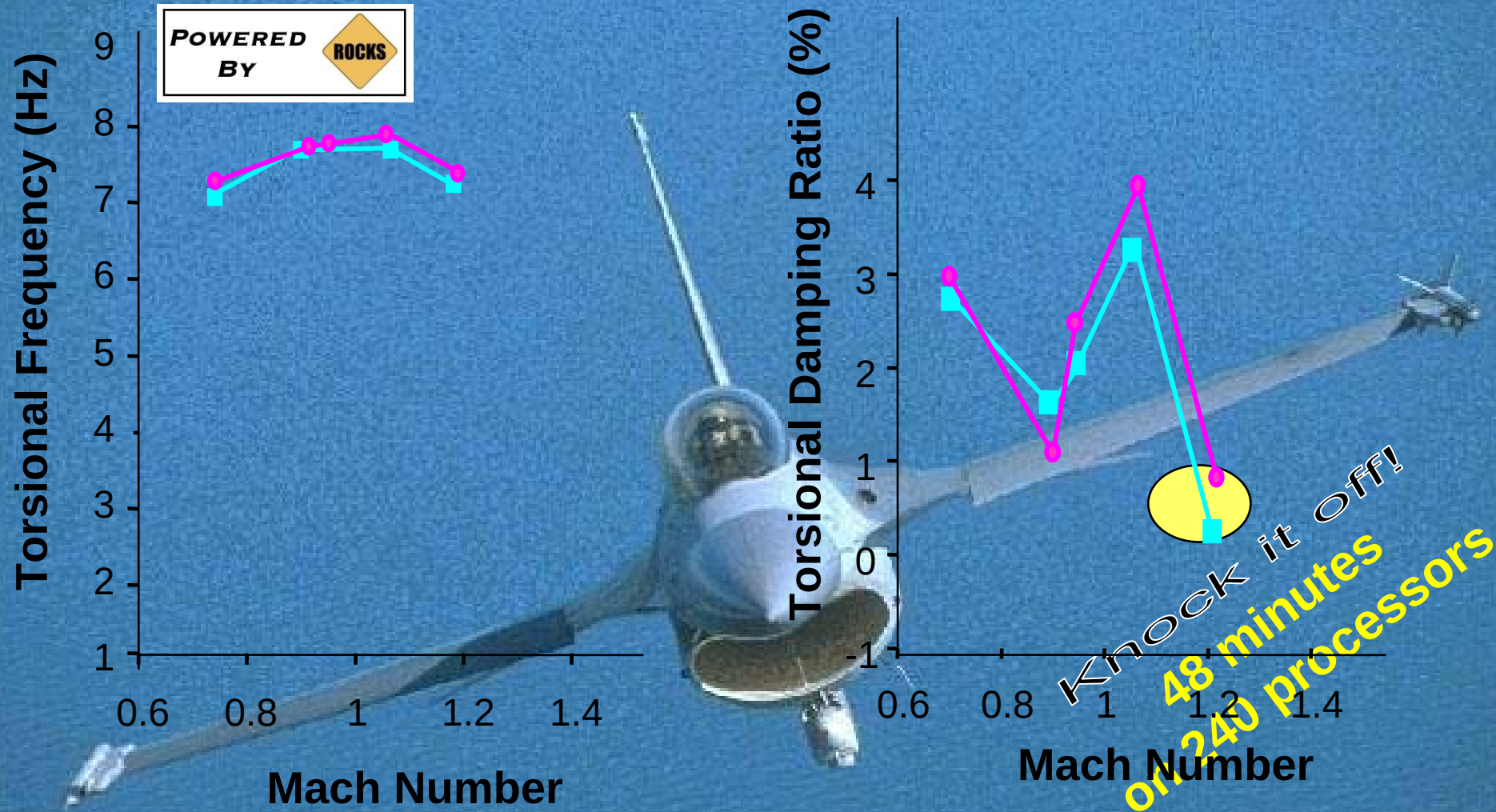
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Fz_stores
Fz_clean

Time : 0.000000

Simulation Versus Flight Data



^ 3D Simulation (Clean Wing)
+ Flight Test Data (Clean Wing)

Databases?



**Desert Storm
(1991)**

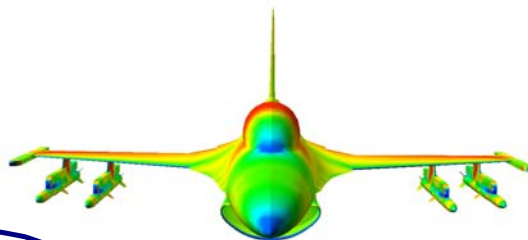
**Iraq War
(2003)**



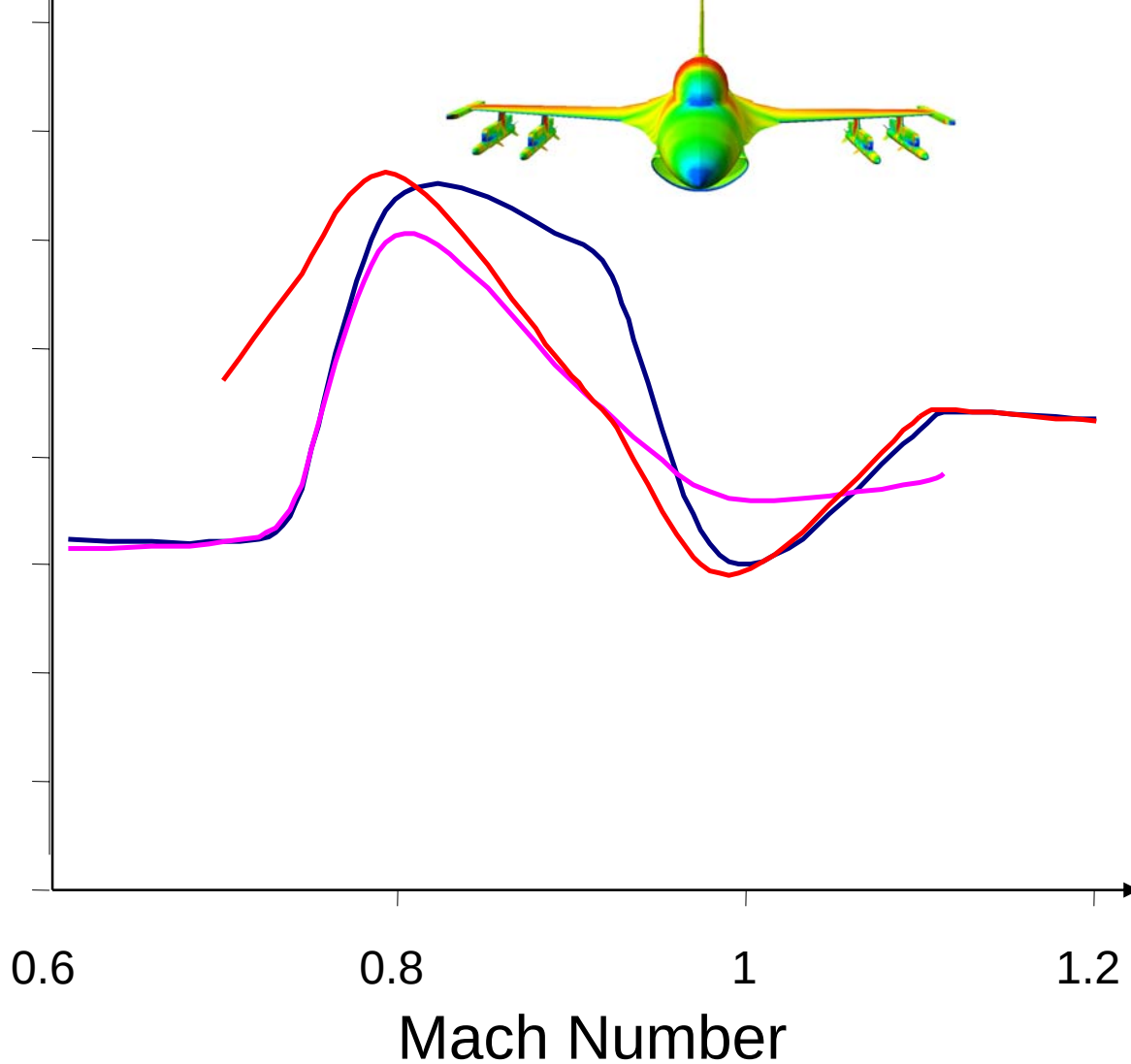
400,000 configurations to be flight tested

Potential

Damping Coefficient (%) -- 1st Torsion



- Flight Test
- FOM (1,170 s.)
- TP-ROM (5 s.)



Nivation



- 164 Processor Intel Xeon 3.0GHz
- 4GB RAM per node
- Myrinet
- Gigabit Ethernet
- Two 1 TB NAS Appliances
- 4 Tools Nodes



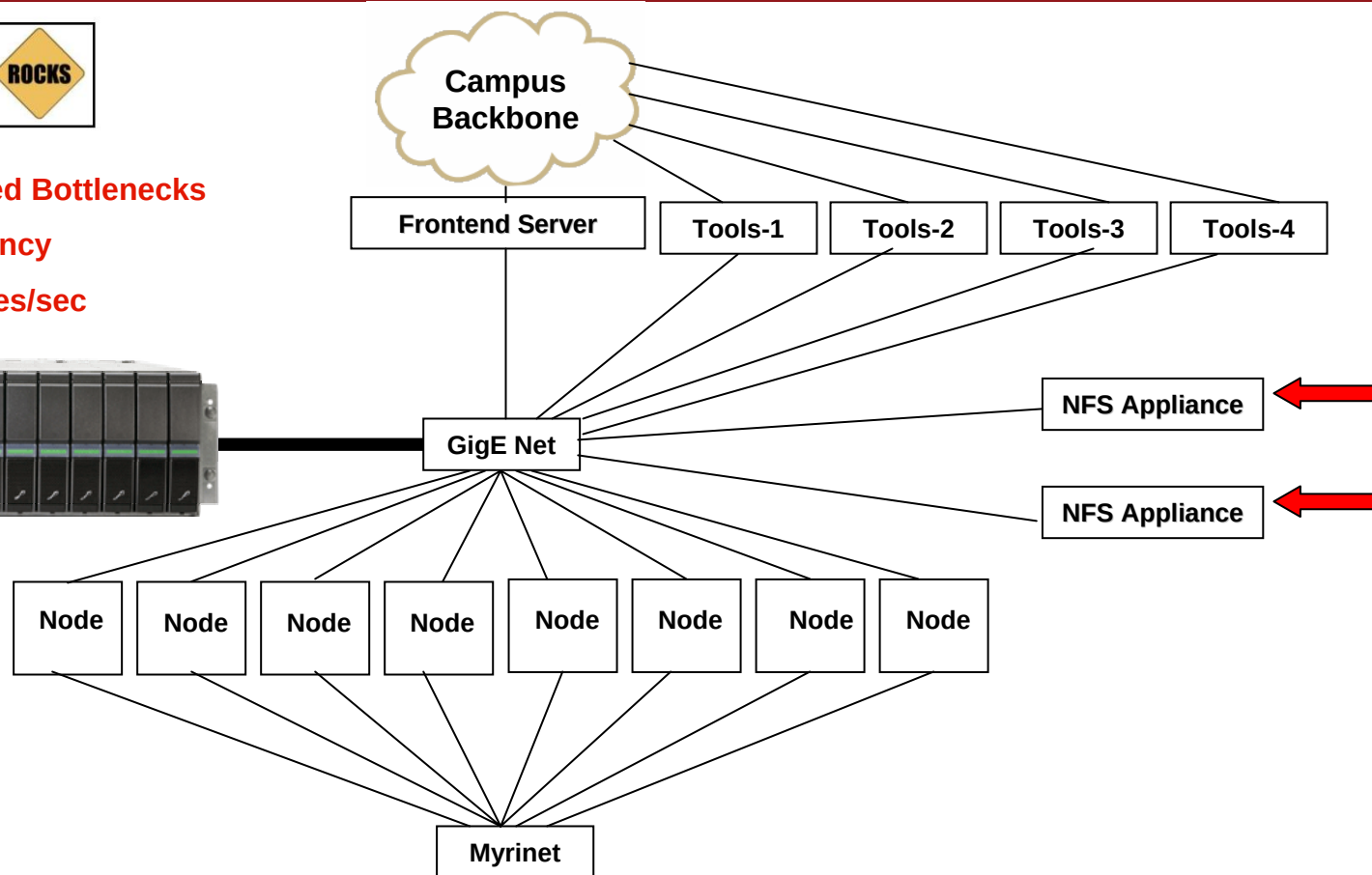
Panasas Eliminates Bottlenecks



Eliminated Bottlenecks

Redundancy

400MBytes/sec



Huge Bottleneck/
Single Point of Failure

Panasas Integration in Less Than 2 Hours



- Installation and configuration of Panasas Shelf - 1 hour
- Switch configuration changes for link aggregation - 10 minutes
- Copy RPM to /home/install/contrib/enterprise/3/public/i386/RPMS - 1 minute
- create/edit extend-compute.xml - 5 minutes

```
# Add panfs to fstab
REALM=10.10.10.10
mount_flags="rw,noauto,panauto"
/bin/rm -f /etc/fstab.bak.panfs
/bin/rm -f /etc/fstab.panfs
/bin/cp /etc/fstab /etc/fstab.bak.panfs
/bin/grep -v "panfs://" /etc/fstab > /etc/fstab.panfs
/bin/echo "panfs://$REALM:global /panfs panfs $mount_flags 0 0" >> /etc/fstab.panfs
/bin/mv -f /etc/fstab.panfs /etc/fstab
/bin/sync

/sbin/chkconfig --add panfs
/usr/local/sbin/check_panfs
LOCATECRON=/etc/cron.daily/slocate.cron
LOCATE=/etc/sysconfig/locate
LOCTEMP=/tmp/slocate.new

/bin/cat $LOCATECRON | sed "s/,proc,/,proc,panfs,/g" > $LOCTEMP
/bin/mv -f $LOCTEMP $LOCATECRON
/bin/cat $LOCATECRON | sed "s/\afs,\afs,\panfs,/g" > $LOCTEMP
/bin/mv -f $LOCTEMP $LOCATECRON
```

- [root@rockcluster]# rocks-dist dist ; cluster-fork '/boot/kickstart/cluster-kickstart' - 30 minutes
- /etc/auto.home userX -fstype=panfs panfs://10.x.x.x/home/userX - script it to save time

Benchmarking Panasas Using Bonnie++



```
#!/bin/bash
```

```
#PBS -N BONNIE
```

```
#PBS -e Log.d/BONNIE.panfs.err
```

```
#PBS -o Log.d/BONNIE.panfs.out
```

```
#PBS -m aeb
```

```
#PBS -M hpcclusters@gmail.com
```

```
#PBS -l nodes=1:ppn=2
```

```
#PBS -l walltime=30:00:00
```

```
PBS_O_WORKDIR='/home/sjones/benchmarks'  
export PBS_O_WORKDIR
```

```
### -----  
### BEGINNING OF EXECUTION  
### -----
```

```
echo The master node of this job is `hostname`  
echo The job started at `date`  
echo The working directory is `echo $PBS_O_WORKDIR`  
echo This job runs on the following nodes:  
echo `cat $PBS_NODEFILE`
```

```
### end of information preamble
```

```
cd $PBS_O_WORKDIR  
cmd="/home/tools/bonnie++/sbin/bonnie++ -s 8000 -n 0 -f -d /home/sjones/bonnie"  
echo "running bonnie++ with: $cmd in directory "`pwd`"  
$cmd >& $PBS_O_WORKDIR/Log.d/run9/log.bonnie.panfs.$PBS_JOBID
```

NFS – 8 Nodes



Version 1.03		Sequential Output				Sequential Input-		Random Seeks	
		-Block-		-Rewrite-		-Block-			
Machine	Size	K/sec	%CP	K/sec	%CP	K/sec	%CP	K/sec	%CP
compute-3-82	8000M	2323	0	348	0	5119	1	51.3	0
compute-3-81	8000M	2333	0	348	0	5063	1	51.3	0
compute-3-80	8000M	2339	0	349	0	4514	1	52.0	0
compute-3-79	8000M	2204	0	349	0	4740	1	99.8	0
compute-3-78	8000M	2285	0	354	0	3974	0	67.9	0
compute-3-77	8000M	2192	0	350	0	5282	0	46.8	0
compute-3-74	8000M	2292	0	349	0	5112	1	45.4	0
compute-3-73	8000M	2309	0	358	0	4053	0	64.6	0

17.80MB/sec for concurrent write using NFS with 8 dual processor jobs

36.97MB/sec during read process

PanFS – 8 Nodes



Version 1.03		Sequential Output				Sequential Input-		Random Seeks	
		-Block-		-Rewrite-		-Block-			
Machine	Size	K/sec	%CP	K/sec	%CP	K/sec	%CP	K/sec	%CP
compute-1-18.	8000M	20767	8	4154	3	24460	7	72.8	0
compute-1-17.	8000M	19755	7	4009	3	24588	7	116.5	0
compute-1-16.	8000M	19774	7	4100	3	23597	7	96.4	0
compute-1-15.	8000M	19716	7	3878	3	25384	8	213.6	1
compute-1-14.	8000M	19674	7	4216	3	24495	7	72.8	0
compute-1-13.	8000M	19496	7	4236	3	24238	7	71.0	0
compute-1-12.	8000M	19579	7	4117	3	23731	7	97.1	0
compute-1-11.	8000M	19688	7	4038	3	24195	8	117.7	0

154MB/sec for concurrent write using PanFS with 8 dual processor jobs

190MB/sec during read process

NFS – 16 Nodes



Version 1.03		Sequential Output				Sequential Input-		Random Seeks	
		-Block-	-Rewrite-			-Block-			
Machine	Size	K/sec	%CP	K/sec	%CP	K/sec	%CP	K/sec	%CP
compute-3-82	8000M	1403	0	127	0	2210	0	274.0	2
compute-3-81	8000M	1395	0	132	0	1484	0	72.1	0
compute-3-80	8000M	1436	0	135	0	1342	0	49.3	0
compute-3-79	8000M	1461	0	135	0	1330	0	53.7	0
compute-3-78	8000M	1358	0	135	0	1291	0	54.7	0
compute-3-77	8000M	1388	0	127	0	2417	0	45.5	0
compute-3-74	8000M	1284	0	133	0	1608	0	71.9	0
compute-3-73	8000M	1368	0	128	0	2055	0	54.2	0
compute-3-54	8000M	1295	0	131	0	1650	0	47.4	0
compute-2-53	8000M	1031	0	176	0	737	0	18.3	0
compute-2-52	8000M	1292	0	128	0	2124	0	104.1	0
compute-2-51	8000M	1307	0	129	0	2115	0	48.1	0
compute-2-50	8000M	1281	0	130	0	1988	0	92.2	1
compute-2-49	8000M	1240	0	135	0	1488	0	54.3	0
compute-2-47	8000M	1273	0	128	0	2446	0	52.7	0
compute-2-46	8000M	1282	0	131	0	1787	0	52.9	0

20.59MB/sec for concurrent write using NFS with 16 dual processor jobs
27.41MB/sec during read process

PanFS – 16 Nodes



Version 1.03		Sequential Output				Sequential Input-		Random Seeks	
		-Block-		-Rewrite-		-Block-			
Machine	Size	K/sec	%CP	K/sec	%CP	K/sec	%CP	K/sec	%CP
compute-1-26	8000M	14330	5	3392	2	28129	9	54.1	0
compute-1-25	8000M	14603	5	3294	2	30990	9	60.3	0
compute-1-24	8000M	14414	5	3367	2	28834	9	55.1	0
compute-1-23	8000M	9488	3	2864	2	17373	5	121.4	0
compute-1-22	8000M	8991	3	2814	2	21843	7	116.5	0
compute-1-21	8000M	9152	3	2881	2	20882	6	80.6	0
compute-1-20	8000M	9199	3	2865	2	20783	6	85.2	0
compute-1-19	8000M	14593	5	3330	2	29275	9	61.0	0
compute-1-18	8000M	9973	3	2797	2	18153	5	121.6	0
compute-1-17	8000M	9439	3	2879	2	22270	7	64.9	0
compute-1-16	8000M	9307	3	2834	2	21150	6	99.1	0
compute-1-15	8000M	9774	3	2835	2	20726	6	77.1	0
compute-1-14	8000M	15097	5	3259	2	32705	10	60.6	0
compute-1-13	8000M	14453	5	2907	2	36321	11	126.0	0
compute-1-12	8000M	14512	5	3301	2	32841	10	60.4	0
compute-1-11	8000M	14558	5	3256	2	33096	10	62.2	0

187MB/sec for concurrent write using PanFS with 8 dual processor jobs

405MB/sec during read process

Capacity imbalances on jobs - 33MB/sec increase from 8 to 16 job run

Panasas Statistics During Write Process



[pancli] sysstat storage

IP	CPU Util	Disk Util	Ops/s In	KB/s Out	Total	Capacity (GB) Avail	Reserved	
10.10.10.250	55%	22%	127	22847	272	485	367	48
10.10.10.253	60%	24%	140	25672	324	485	365	48
10.10.10.245	53%	21%	126	22319	261	485	365	48
10.10.10.246	55%	22%	124	22303	239	485	366	48
10.10.10.248	57%	22%	134	24175	250	485	369	48
10.10.10.247	52%	21%	124	22711	233	485	366	48
10.10.10.249	57%	23%	135	24092	297	485	367	48
10.10.10.251	52%	21%	119	21435	214	485	366	48
10.10.10.254	53%	21%	119	21904	231	485	367	48
10.10.10.252	58%	24%	137	24753	300	485	366	48
Total "Set 1"	55%	22%	1285	232211	2621	4850	3664	480

Sustained BW 226 MBytes/Sec during 16 1GB concurrent writes

Panasas Statistics During Read Process



[pancli] sysstat storage

IP	CPU		Disk		Ops/s		KB/s		Capacity (GB)	
	Util	Util	In	Out	Total	Avail	Reserved			
10.10.10.250	58%		95%	279	734	21325	485	355	48	
10.10.10.253	60%		95%	290	727	22417	485	353	48	
10.10.10.245	54%		92%	269	779	19281	485	353	48	
10.10.10.246	59%		95%	290	779	21686	485	354	48	
10.10.10.248	60%		95%	287	729	22301	485	357	48	
10.10.10.247	52%		91%	256	695	19241	485	356	48	
10.10.10.249	57%		93%	276	708	21177	485	356	48	
10.10.10.251	49%		83%	238	650	18043	485	355	48	
10.10.10.254	45%		82%	230	815	15225	485	355	48	
10.10.10.252	57%		94%	268	604	21535	485	354	48	
Total "Set 1"	55%		91%	2683	7220	202231	4850	3548	480	

Sustained BW 197 MBytes/Sec during 16 1GB concurrent sequential reads

Typical Storage Utilization with Cluster at 76%



[pancli] sysstat storage

IP	CPU	Disk	Ops/s		KB/s	Capacity (GB)		
	Util	Util	In	Out	Total	Avail	Reserved	
10.10.10.250	6%	5%	35	292	409	485	370	48
10.10.10.253	5%	4%	35	376	528	485	368	48
10.10.10.245	4%	3%	29	250	343	485	368	48
10.10.10.246	6%	4%	28	262	373	485	369	48
10.10.10.248	5%	3%	27	234	290	485	372	48
10.10.10.247	3%	3%	1	1	2	485	370	48
10.10.10.249	5%	3%	48	258	365	485	371	48
10.10.10.251	4%	3%	46	216	267	485	369	48
10.10.10.254	4%	3%	32	256	349	485	370	48
10.10.10.252	4%	3%	34	337	499	485	370	48
Total	4%	3%	315	2482	3425	4850	3697	480

sustained BW 2.42 Mbytes/sec in - 3.34 Mbytes/sec out

[root@frontend-0 root]# showq

ACTIVE JOBS-----

JOBNAME	USERNAME	STATE	PROC	REMAINING	STARTTIME
8649	sjones	Running	2	1:04:55:08	Sun May 15 18:20:23
8660	user1	Running	6	1:23:33:15	Sun May 15 18:58:30
8524	user2	Running	16	2:01:09:51	Fri May 13 20:35:06
8527	user3	Running	16	2:01:23:19	Fri May 13 20:48:34
8590	user4	Running	64	3:16:42:50	Sun May 15 10:08:05
8656	user5	Running	16	4:00:55:36	Sun May 15 18:20:51
8647	user6	Running	5	99:22:50:42	Sun May 15 18:15:58

7 Active Jobs 125 of 164 Processors Active (76.22%)

65 of 82 Nodes Active (79.27%)

Conclusion



- Summary
- Demo
- Questions and Answers



Thank you

For more information about Steve Jones:
High Performance Computing Clusters

<http://www.hpcclusters.org>

For more information about Panasas:

<http://www.panasas.com>

For more information about Rocks:

<http://www.rocksclusters.org>