

# AI Big-Model and OR Mathematical Optimization

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# Today's Sharing

- 1. AI Big-Model and OR/Math Optimization**
- 2. Online Training Algorithms and Applications**
- 3. Pre-Trained Mixed Integer Linear Programming Solvers and Applications**
- 4. Big-Model Training Needs Efficient Optimization Algorithms?**

# ChatGPT的能力边界 | 求解数学优化问题?

- ▶ 优化问题根据变量、约束和目标函数的特性，可以分为不同种类，如线性规划、非线性规划和整数规划等
- ▶ 小规模整数优化问题，可以通过穷举法求解
- ▶ 大规模整数规划问题一般由专业的整数规划软件求解 (Gurobi, COPT)

## 一个简单的数学规划问题

- ▶ 假设有一个最多能承重5kg的背包和5个重量不同的物品，每个物品都具有一定的价值。将哪些物品装入背包，可以使得背包内的物品总价值最大？
- ▶ 物品的重量和价值如下表所示：

物品	1	2	3	4	5
重量 (kg)	2	1	2.5	3.5	1.5
价值 (元)	18	10	20	28	15



# 使用ChatGPT求解科学决策问题

## 将上述背包问题向ChatGPT提问：



假设有5件物品和一个最多能承重5kg的背包。第1件物品的重量是2kg，其价值为18元；第2件物品的重量是1kg，其价值为10元；第3件物品的重量是2.5kg，其价值为20元；第4件物品的重量是3.5kg，其价值为28元；第5件物品的重量是1.5kg，其价值为15元。请问将哪些物品装入背包里，物品的价值总和最大？请使用整数规划求解该问题。



这是一个经典的背包问题，可以使用整数规划方法求解。

我们可以定义0/1变量 $x_i$ ，表示是否将第 $i$ 件物品放入背包中。同时，我们还需要定义一个非负整数变量 $z$ ，表示物品的总价值。

然后，我们可以将问题转化为以下的线性规划问题：

$$\text{最大化 } z = 18x_1 + 10x_2 + 20x_3 + 28x_4 + 15x_5$$

满足约束条件：

$$2x_1 + x_2 + 2.5x_3 + 3.5x_4 + 1.5x_5 \leq 5$$

$$x_i \in \{0,1\}, i=1,2,3,4,5$$

由于 $x_i$ 是整数变量，这是一个整数规划问题。可以使用整数规划求解器求解。

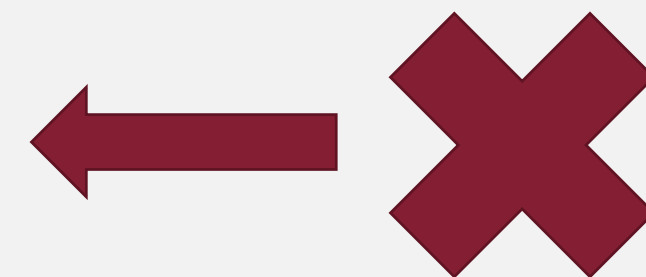
求解结果表明，将第1件、第3件、第5件物品装入背包中，可以使得物品的总价值最大，为 $18+20+15=53$ 元。



ChatGPT 在不调用求解器的前提下，能否求解大规模优化问题？



在不调用求解器的前提下，对于大规模的优化问题，通常也比较难直接求解得到最优解，因为现实中的大多数优化问题都非常复杂，存在着多个目标、多个约束、非线性、离散等多种问题，很容易遇到组合爆炸的问题。

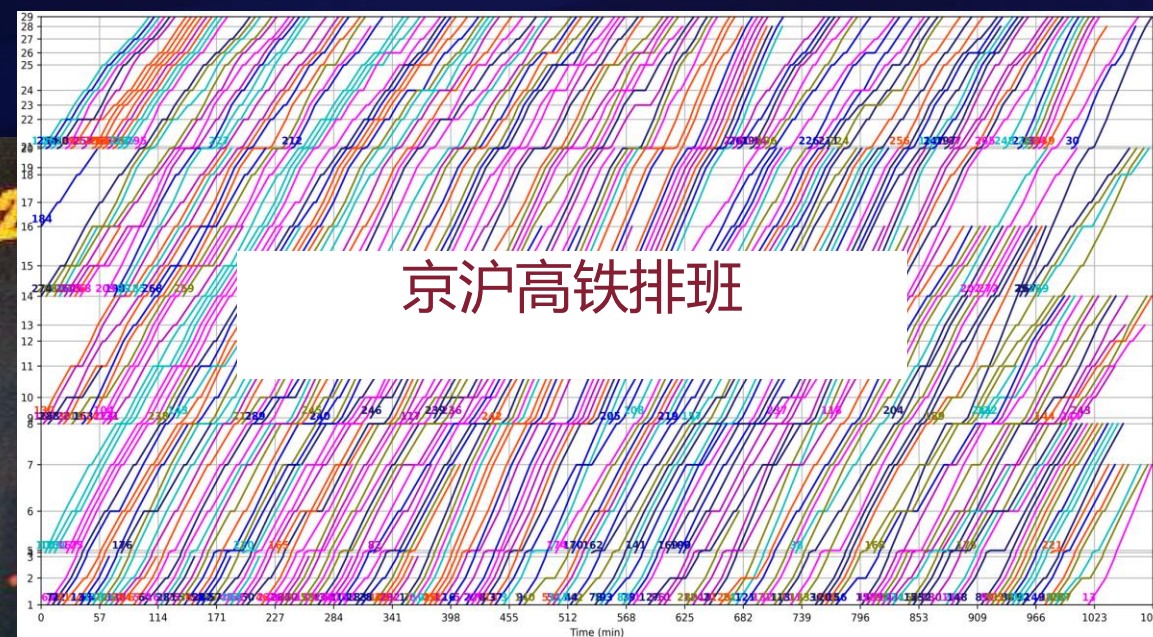
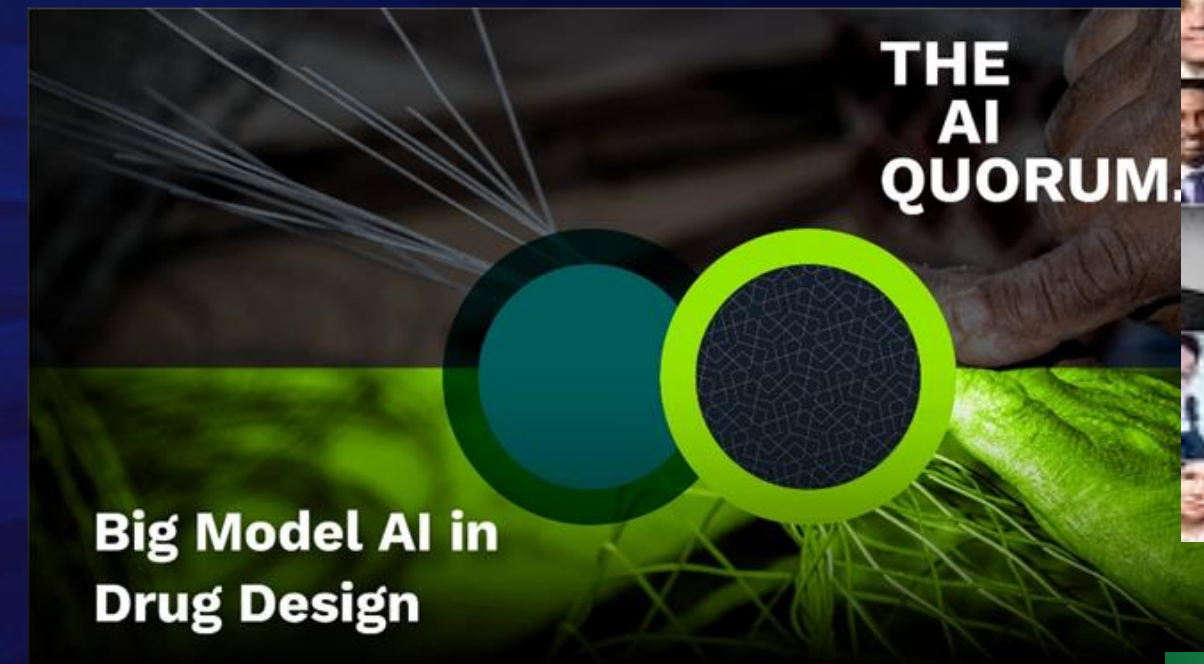
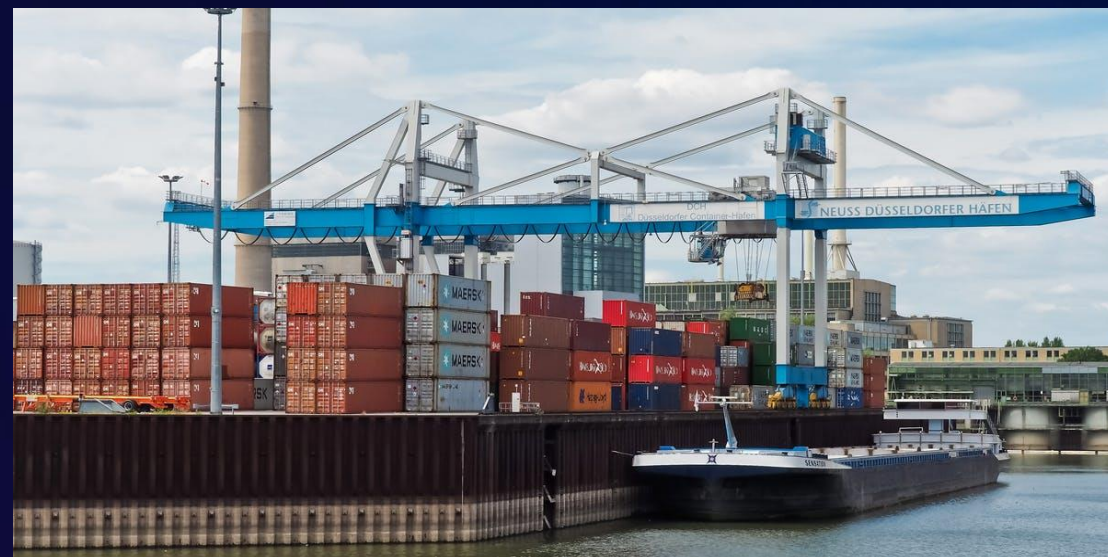
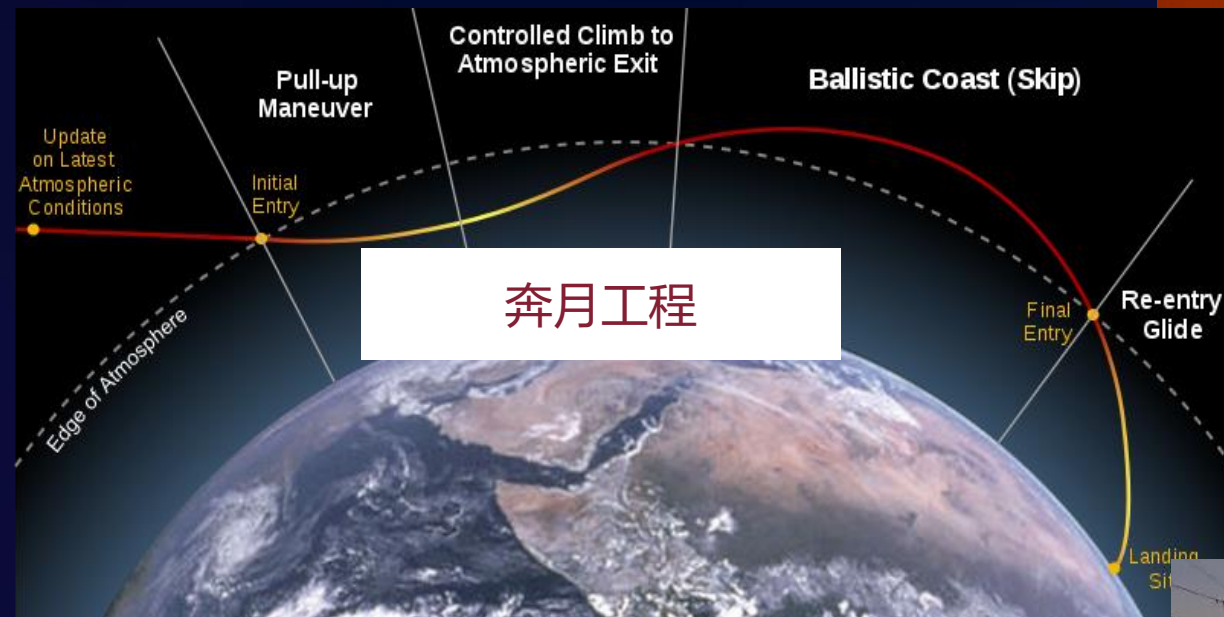


## 向ChatGPT提问：

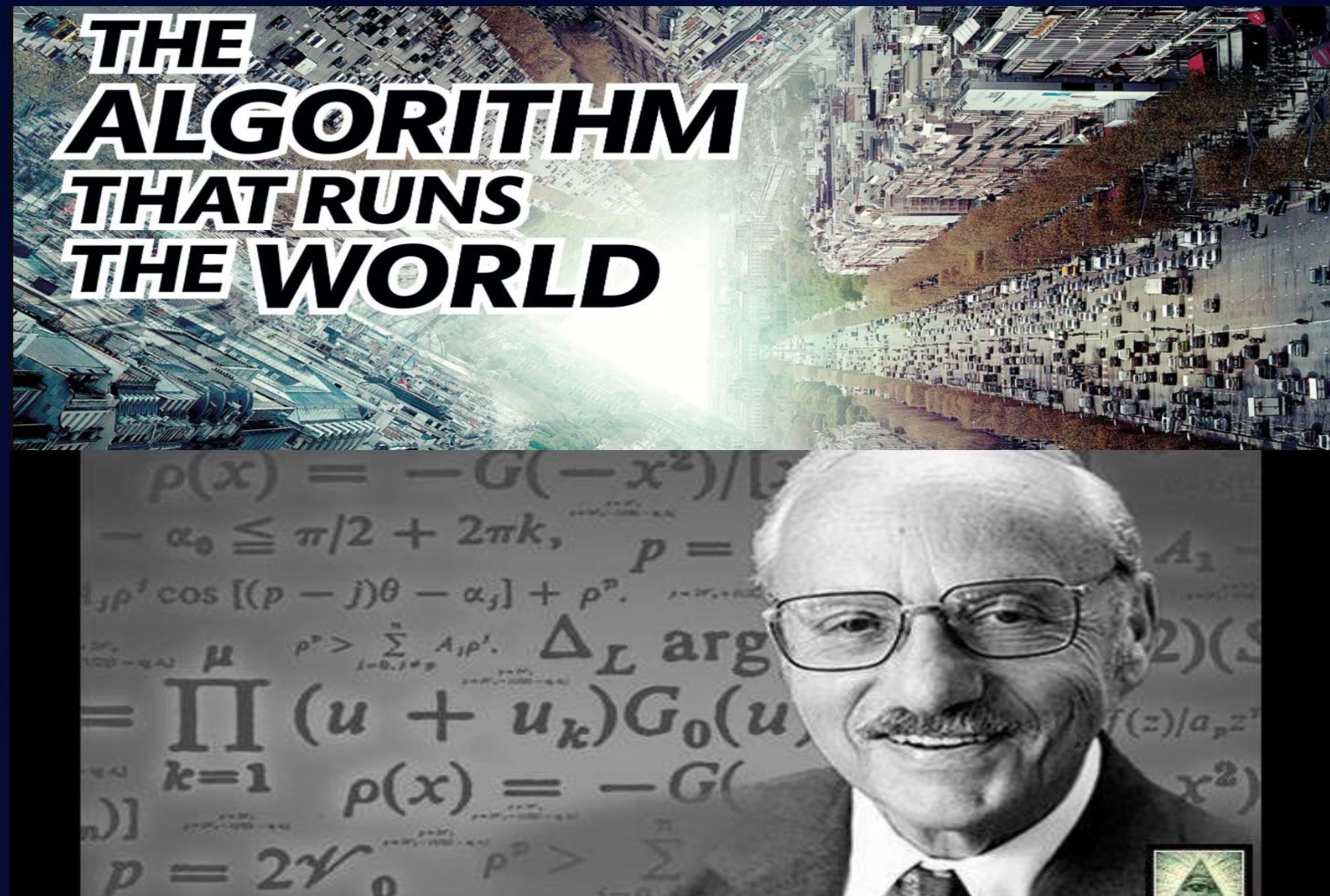
- ChatGPT 会尝试性给出解答，但是不能保证答案的正确性。
- ChatGPT会建议调用求解器求解优化问题：全科vs专科
- ChatGPT 并不是万能的，它是一个语言处理模型，回答复杂的计算问题还需要调用专业化的软件。

上述问答产生于2023年3月12日

# Applications of OR and AI Models



# Differences of OR and AI Models



- Based on Science/Logic
- Physical/Economical Principles
  - Objective
  - Definitive
- Explainable Insights
- Online Training&Decision-Making

- Based on Cases/Experience
  - Observation/Behaviour
    - Subjective
    - Probabilistic
    - Black-box
- Offline Training

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# Linear Programming and LP Giants won Nobel Prize...

$$\max \sum \pi_j x_j$$

$$\text{s.t. } \sum_j a_j x_j \leq b,$$

$$0 \leq x_j \leq 1 \quad \forall j = 1, \dots, n$$

$$\min \mathbf{b}^T \mathbf{p} + \sum \max\{0, \pi_j - \mathbf{a}_j^T \mathbf{p}\}$$

$$\text{s.t. } \mathbf{p} \geq \mathbf{0}$$





# Online Linear Programming: an Online Auction Example

- There is a fixed selling period or number of buyers; and there is a fixed inventory of goods
- Customers come and require a bundle of goods and make a bid
- Decision: To sell or not to sell to each individual customer on the fly?
- Objective: Maximize the revenue.

Bid #	\$100	\$30	....	...	...	Inventory
Decision	x1	x2				
Pants	1	0	....	...	...	100
Shoes	1	0				50
T-Shirts	0	1				500
Jackets	0	0				200
Hats	1	1	...	...	...	1000

# Online Linear Programming Model and Theory

- OLP theory and practice (Agrawal et al. 2010, 14, Li&Y 2022)

$$\begin{aligned} \max \quad & \sum \pi_j x_j \\ \text{s.t.} \quad & \sum_j a_j x_j \leq b, \\ & 0 \leq x_j \leq 1 \quad \forall j = 1, \dots, n \end{aligned}$$

- OLP:

- Variables together with their data points arrive sequentially and decision makers need decide  $x_j$  on the fly, that is, before knowing the “future” data points
- Learning-while-Doing vs Learning-First and Deciding-Second (collect and learn all relevant data, then solve for all  $\mathbf{x}$ )
- Offline LP’s objective value is a upper bond for the online version
- Is there an optimal online decision algorithm/mechanism<sup>10</sup>

# Price Mechanism for Online Auction

- Learn and compute itemized optimal prices
- Use the prices to price each bid internally
- Accept if it is a over bid, and reject otherwise
- There is an Optimal Online Algorithm to achieve the best you could do!

Massive episodes are transferred into Knowledges that can be stored/reused

Bid #	\$100	\$30	....	...	...	Inventory	Price?
Decision	x1	x2					
Pants	1	0	....	...	...	100	45
Shoes	1	0				50	45
T-Shirts	0	1				500	10
Jackets	0	0				200	55
Hats	1	1	...	...	...	1000	15



$$\max \sum \pi_j x_j \quad \text{s.t.} \quad \sum_j a_j x_j \leq \mathbf{b}, \quad x_j \geq 0 \quad \forall j = 1, \dots, J$$

- The decision variable  $x_j$  represents the **total-times of pulling** the  $j$ -th arm.
- We have developed a two-phase algorithm
  - **Phase I**: Distinguish the optimal **super-basic** variables/arms from the optimal **non-basic** variables/arms with as fewer number of plays as possible
  - **Phase II**: Use the arms in the optimal face to exhaust the resource through an adaptive procedure and achieve **fairness**
- The algorithm achieves a problem dependent regret that bears a **logarithmic** dependence on the horizon  $T$ . Also, it identifies a number of LP-related parameters as the **bottleneck or condition-numbers** for the problem
  - Minimum non-zero **reduced cost**
  - Minimum **singular-values** of the optimal basis matrix.
- **First algorithm** to achieve the  $O(\log T)$  regret/gap bound [Li, Sun & Y 2021 ICML] (<https://proceedings.mlr.press/v139/li21s.html>)

# 阿里巴巴在2019年云栖大会上提到在智能履行决策上使用OLP的算法

2018 杭州·云栖大会 Alibaba Group

## 智能履行决策

商家

杭州-上海 杭州-广州 杭州-北京 杭州-武汉 ...

YTO ZTO YUNDA

菜鸟智能发货引擎

时效	服务	成本	单量平衡	...
线路容量	网点容量	局部优化	全局优化	...

最优快递

智能决策 ML & Optimization

商家的履行是带有全局约束的序列执行决策

- Online assignment problem
- Control based method
- Online linear programming

Ref: Agrawal, Shipra, Zizhuo Wang, and Yinyu Ye. "A dynamic near-optimal algorithm for online linear programming." *Operations Research* 62.4 (2014): 876-890.

决策变量

$$C_{ij} = c1 * \text{成本} + c2 * \text{服务} + c3 * \text{时效}$$
$$\max_x \sum_{i=1}^n \sum_{j=1}^m C_{ij} x_{ij}$$

将订单 I 匹配给快递公司 j 与否

$$\text{s.t.} \sum_{j=1}^m x_{ij} \leq 1$$
$$\sum_{i=1}^n x_{ij} * a_j \leq u_j$$

商家发货CP总单量比例约束

$$\sum_{i=1}^n \sum_{j=1}^m x_{ij} b_{k,ij} \leq B_k$$

全局约束值, 比如总成本

阿里巴巴团队在2020年CIKM会议论文Online Electronic Coupon Allocation based on Real-Time User Intent Detection上提到他们设计的发红包的机制也使用了OLP的方法 [2]

### Spending Money Wisely: Online Electronic Coupon Allocation based on Real-Time User Intent Detection

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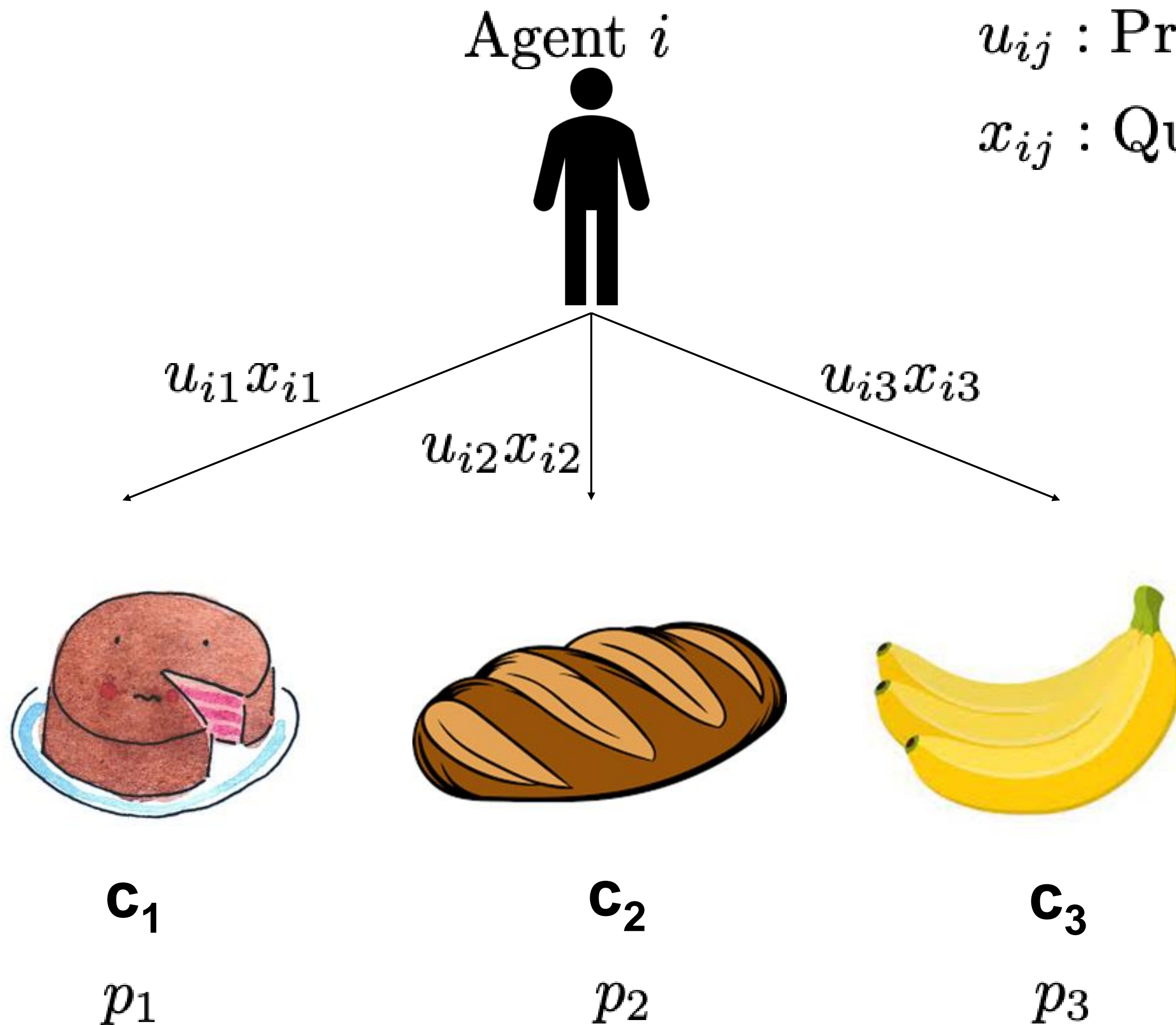
$$\begin{aligned} & \max \sum_{i=1}^M \sum_{j=1}^N v_{ij} x_{ij} \\ & \text{s.t.} \sum_{i=1}^M \sum_{j=1}^N c_j x_{ij} \leq B, \\ & \sum_j x_{ij} \leq 1, \quad \forall i \\ & x_{ij} \geq 0, \quad \forall i, j \end{aligned} \quad (5)$$

### 3.3 MCKP-Allocation

We adopt the primal-dual framework proposed by [2] to solve the problem defined in Equation 5. Let  $\alpha$  and  $\beta_j$  be the associated dual variables respectively. After obtaining the dual variables, we can solve the problem in an online fashion. Precisely, according to the principle of the primal-dual framework, we have the following allocation rule:

$$x_{ij} = \begin{cases} 1, & \text{where } j = \arg \max_i (v_{ij} - \alpha c_j) \\ 0, & \text{otherwise} \end{cases} \quad (9)$$

# A key resource allocation model to achieve efficient allocation is the Fisher Market



$u_{ij}$  : Preference of Agent  $i$  for one unit of good  $j$

$x_{ij}$  : Quantity of good  $j$  purchased by person  $i$

$p_j$  : Price of Good  $j$

$w_i$  : Budget of Agent  $i$

**Individual Optimization Problem:**

$$\max_{\mathbf{x}_i} \sum_j u_{ij} x_{ij}$$

$$\text{s.t. } \mathbf{p}^T \mathbf{x}_i \leq w_i$$

$$\mathbf{x}_i \geq \mathbf{0}$$

**Do Prices exist to CLEAR the market?**  
(Prices are posted and known to all agents so that they have freedom to choose)<sup>6</sup>

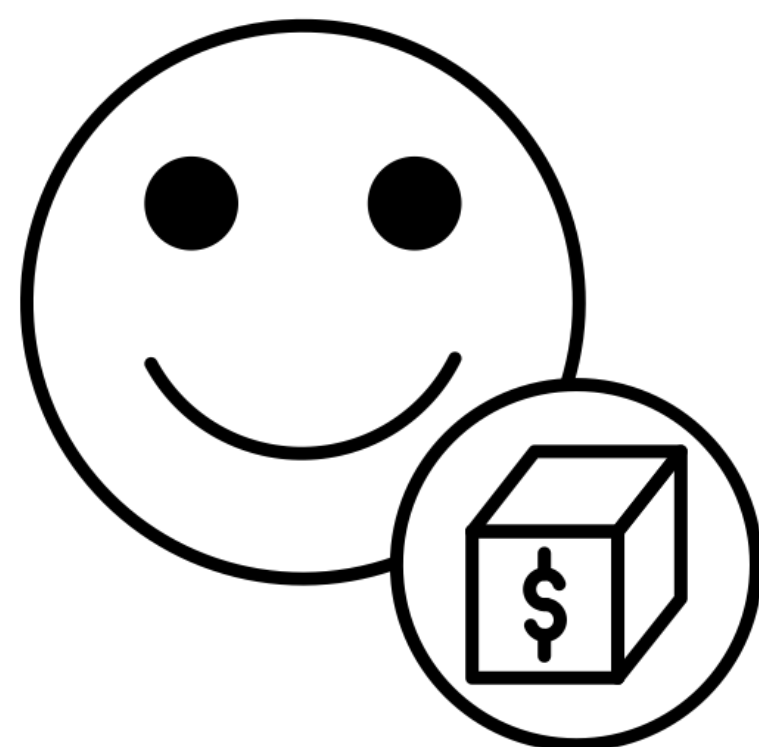
$M =$  Total Number of Goods



# The prices can be derived from a centralized optimization problem with a budget-weighted social objective (Eisenberg-Gale)

**Individual Optimization Problem:**

$$\begin{aligned} \max_{\mathbf{x}_i} \quad & \sum_j u_{ij} x_{ij} \\ \text{s.t.} \quad & \mathbf{p}^T \mathbf{x}_i \leq w_i \\ & \mathbf{x}_i \geq \mathbf{0} \end{aligned}$$



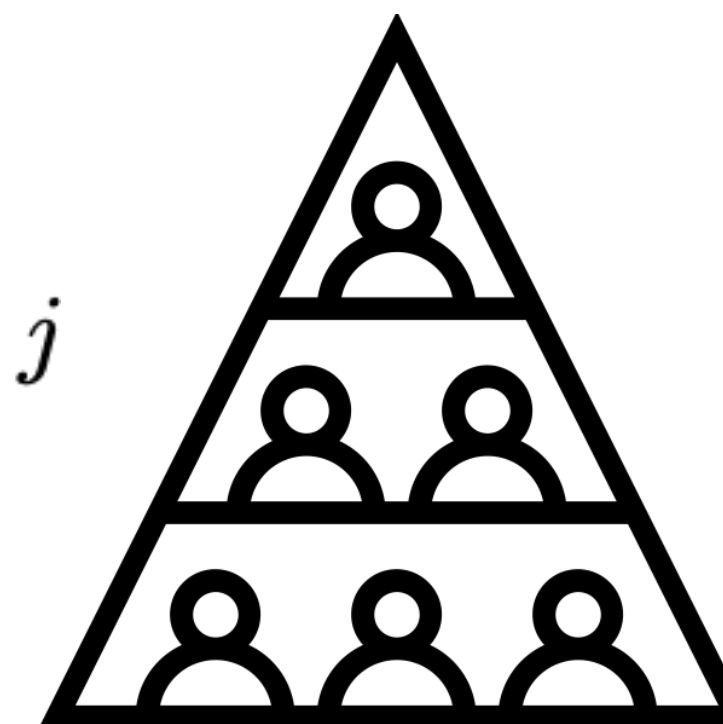
**Social Optimization Problem:**

$$\begin{aligned} \max_{\mathbf{x}_i, \forall i \in [N]} \quad & \sum_i w_i \log \left( \sum_j u_{ij} x_{ij} \right) \\ \text{s.t.} \quad & \sum_i x_{ij} \leq c_j, \forall j \in [M] \\ & x_{ij} \geq 0, \forall i, j \end{aligned}$$

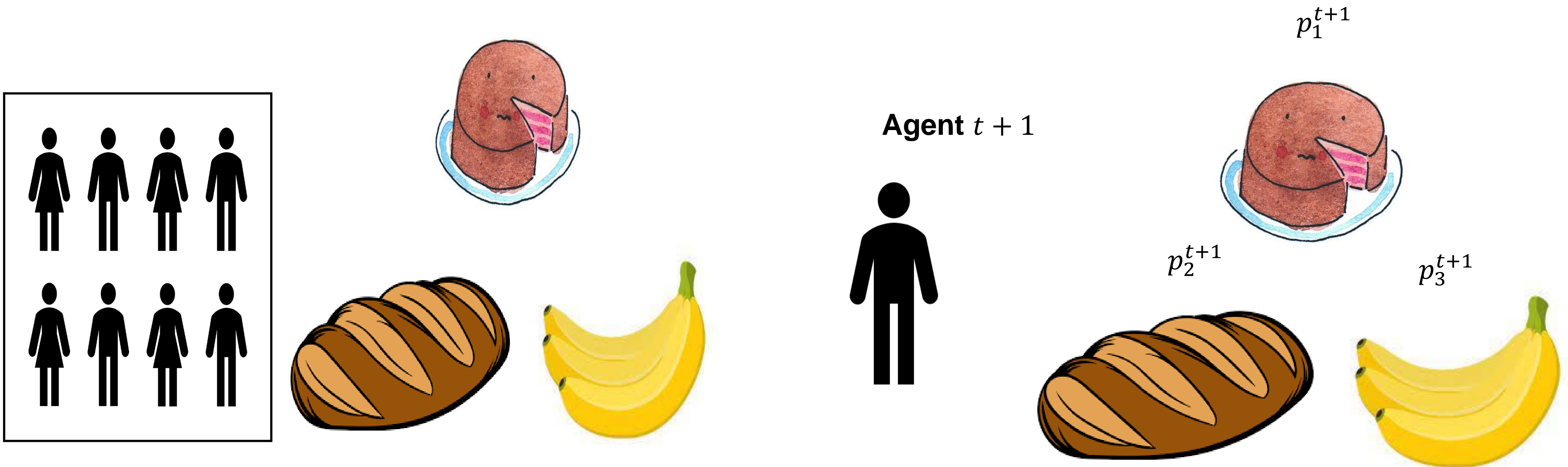
Capacity Constraints

$p_j$  : Price of Good  $j$  = Dual Variable of Constraint  $j$

$C_j$  can be decision variables subject to other resource constraints



# Online Market Pricing: How to update posted-prices to minimize regret of the Eisenberg/Gale social welfare while achieving market clearness



Static Fisher Market Price Equilibrium Theory and Limit  
Now Agents Arrive Online

The price at time  $t + 1$  is updated and reposted based on observed consumption  $x^t$  at time  $t$   
*Jalota and Y*  
<https://arxiv.org/abs/2205.00825>

# We develop a revealed preference algorithm with sub-linear regret and constraint violation guarantees

---

## Algorithm 2: Revealed Preference Algorithm for Online Fisher Markets

---

**Input** : Number of users  $n$ , Vector of good capacities per user  $\mathbf{d} = \frac{\mathbf{c}}{n}$

Initialize  $\mathbf{p}^1 > \mathbf{0}$  ;

**for**  $t = 1, 2, \dots, n$  **do**

**Phase I** ;

    User purchases an optimal bundle of goods  $\mathbf{x}_t$  given the price  $\mathbf{p}^t$  ;

**Phase II (Price Update)** ;

$\mathbf{p}^{t+1} \leftarrow \mathbf{p}^t - \gamma_t (\mathbf{d} - \mathbf{x}_t)$  ;

Difference between market share of each agent and goods purchased

**end**

---

Step-size:  $O\left(\frac{1}{\sqrt{n}}\right)$

Only requires knowledge of user consumption (and not their budgets or utilities) to update prices

**Theorem:** Under i.i.d. budget and utility parameters with strictly positive support and when good capacities are  $O(n)$ , Algorithm 2 achieves an expected regret of  $R_n(\boldsymbol{\pi}) \leq O(\sqrt{n})$  and expected constraint violation of  $V_n(\boldsymbol{\pi}) \leq O(\sqrt{n})$ , where  $n$  is the number of arriving users.

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# MILPI: Unit Commitment Problem

- Electricity is generated from units (**various** generators)
- Transmitted **safely** and **stably** through power grids
- Consumed at **minimum (reasonable)** price



*Optimization has its role to play*

minimize **Cost of electricity**  
subject to **Safety and Stability**

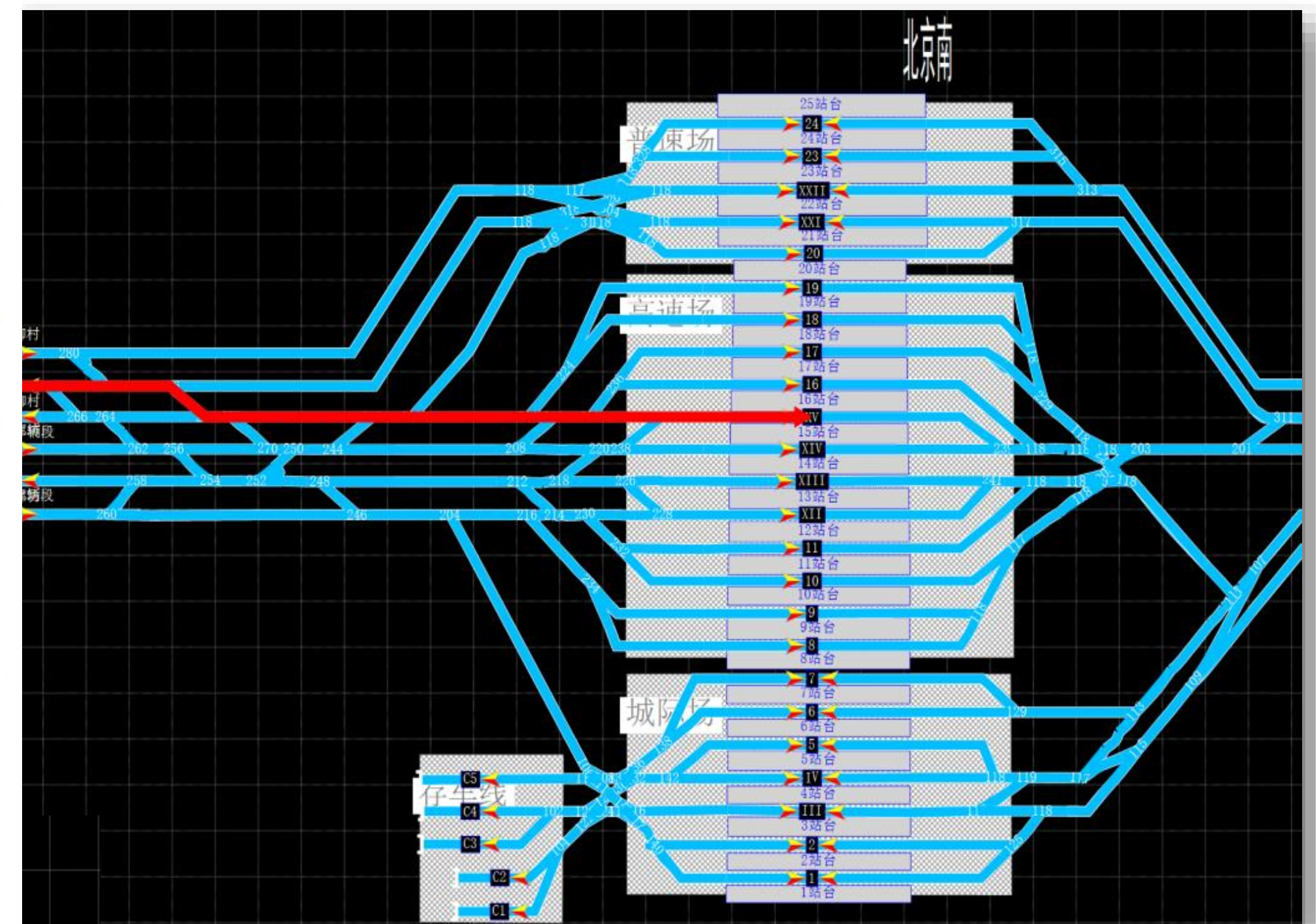
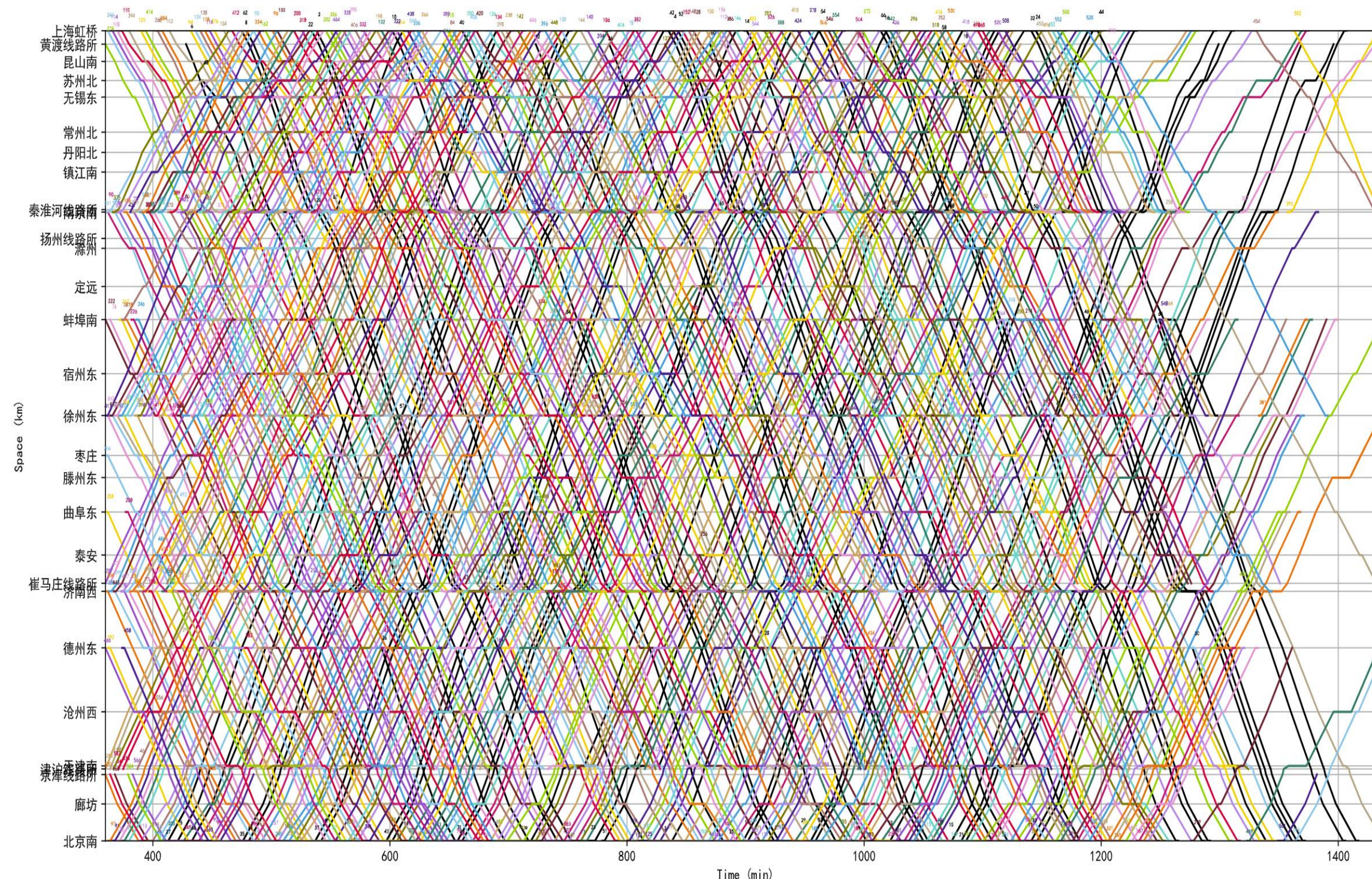
Adaptivity to various units

Unit commitment problem dispatches the units **safely** and **stably** at **minimum** cost



# MILP II: Beijing-Shanghai High-speed Railway Scheduling Optimization

## COPT, Cardinal Operations 2022



# Risk-Pooling Cuts in MILP

- Given an MILP, the fraction solution tells us

$$\begin{pmatrix} \hat{y}_1(\xi) \\ \hat{y}_2(\xi) \\ \vdots \\ \hat{y}_n(\xi) \end{pmatrix} = \begin{pmatrix} 0.99 \\ 0.12 \\ \vdots \\ 0.38 \end{pmatrix}$$

- Each  $\hat{y}$  is the likelihood a variable takes 1 or 0 in the optimal solution
- Each variable introduces some **risk/variance** of **such rounding**  
so that dealing them separately results in **extremely risk outcomes**

Q: What should we do seeing a set of risky guesses?      A: **Put them in a pool!**

# Risk-Pooling Cuts via Moment Ambiguity of DRO

- Pooling the binary variables by adding “confidence” cardinality cuts

$$\sum_{i \in \mathcal{U} = \{j: \hat{y}_j(\xi) \geq 0.9\}} y_i^*(\xi) \geq \alpha \cdot |\mathcal{U}| \quad \sum_{i \in \{\mathcal{L} = \{j: \hat{y}_j(\xi) \leq 0.1\}\}} y_i^*(\xi) \leq \beta \cdot |\mathcal{L}|$$

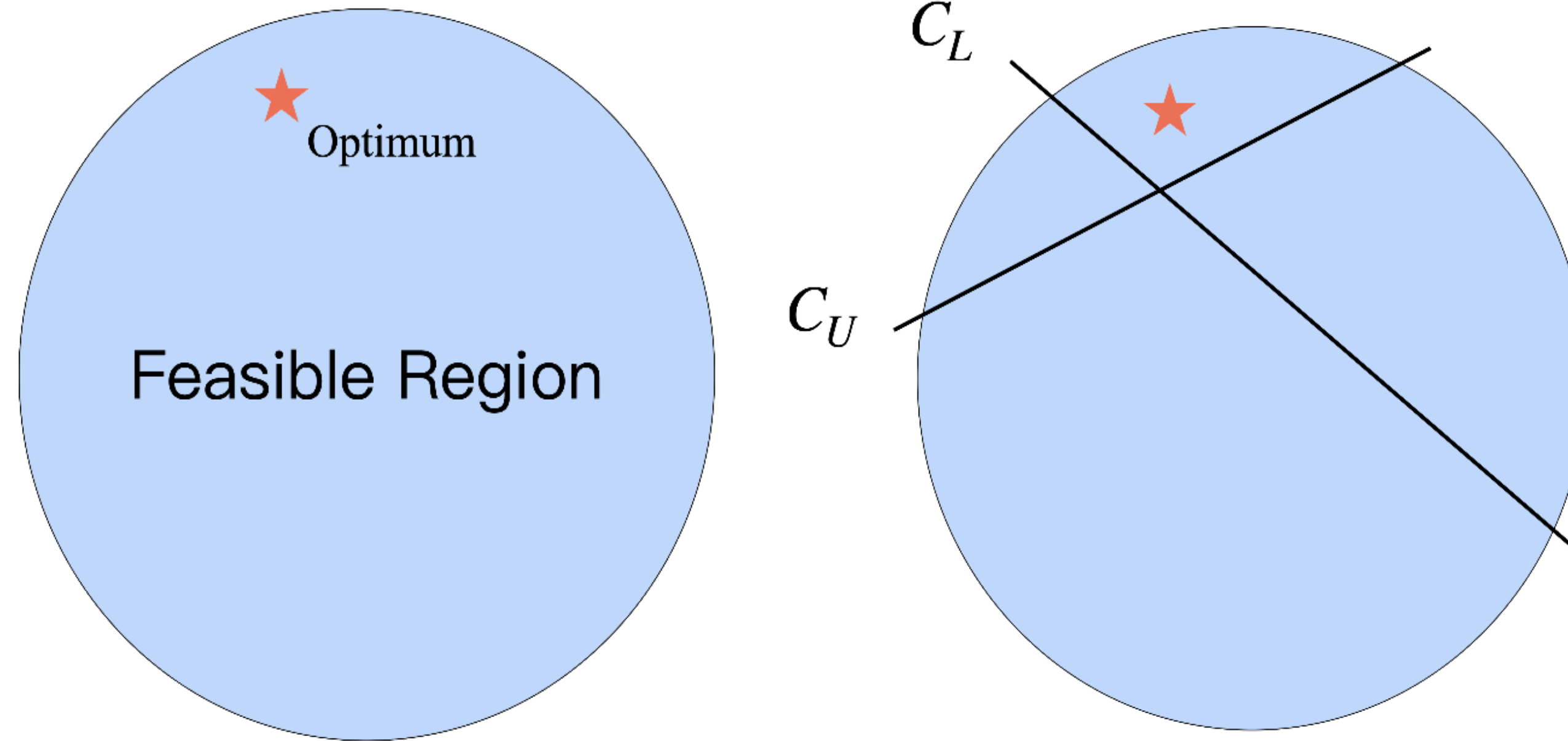
- Intuitively we know that the above two inequalities are **expectedly** to hold for  $\alpha \rightarrow 0.9$  and  $\beta \rightarrow 0.1$
- These two inequalities are exactly *cutting planes* for MILP
- Choose  $\alpha, \beta$  **to increase the confidence level:**

Interpret  $y_j^*$  as some random variables with expectation  $\hat{y}_j$ , then justify it

**Massive data are transferred into Knowledges that can be stored/reused**



# Statistical Confidence Cut Generation (Gao et al. SHUFE, 2023)



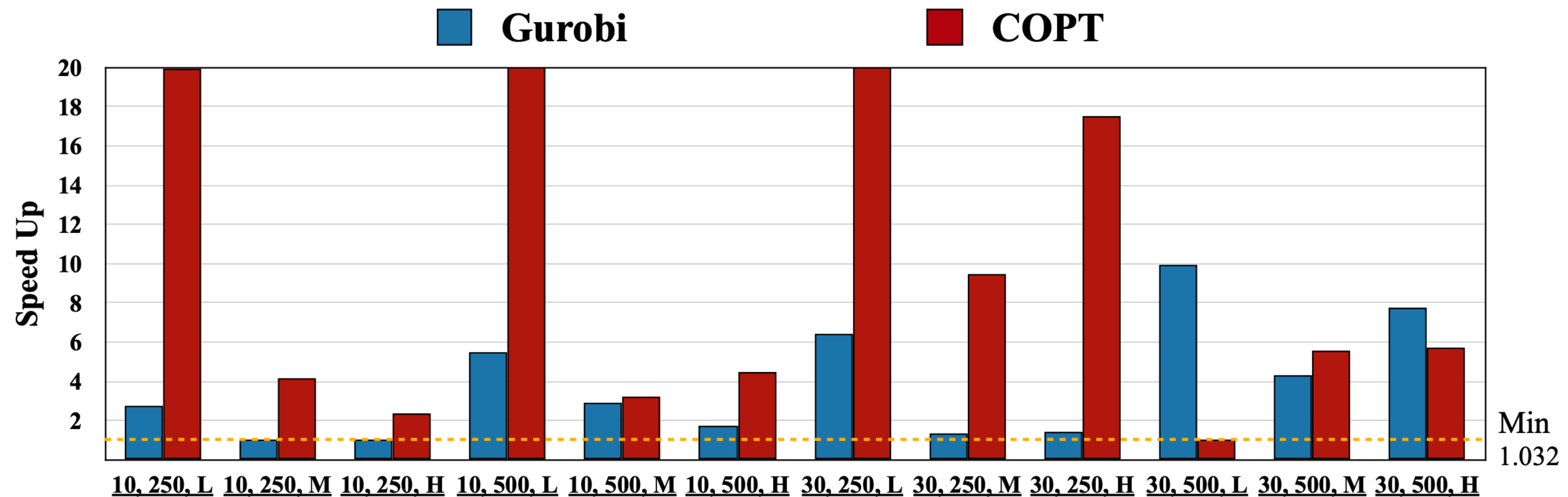
- Overall, the two cuts (and their complement) split the whole feasible region into four regions
- Solving the **most likelihood** region of two cuts often gives a satisfying solution with **confidence**
- Branching over all four regions **independently** will not miss the optimal solution

# Numerical Test Results I

- The method is tested on multi-knapsack, set-covering and unit-commitment problems

Train from 500 instances and test on 20 instances

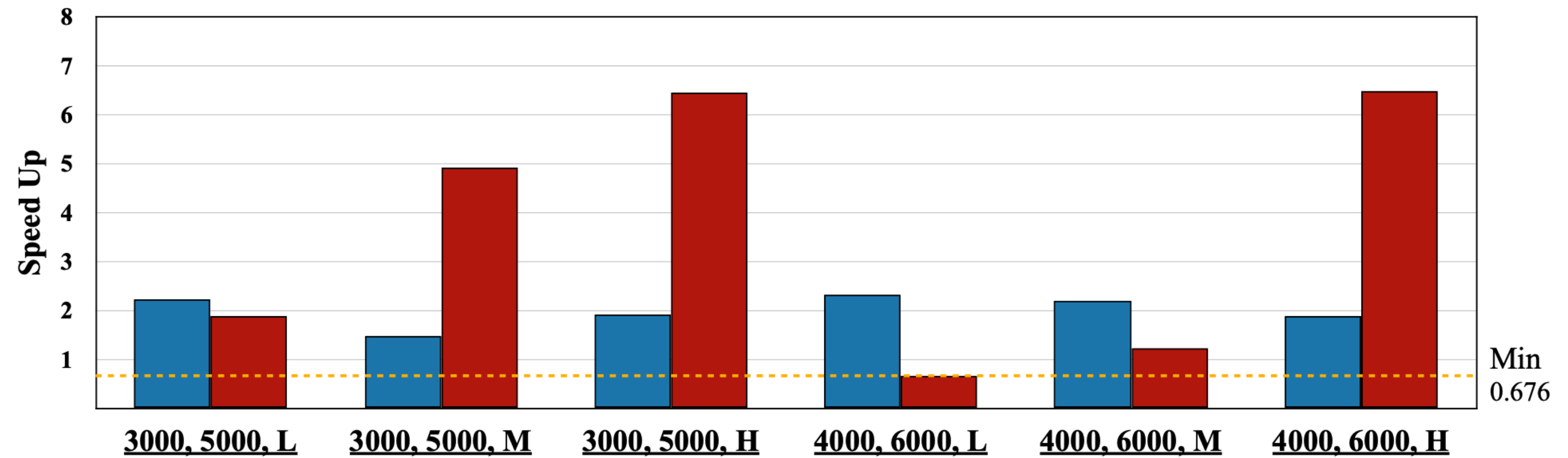
- Measure the speedup of finding a good solution on in the region formed by two cuts



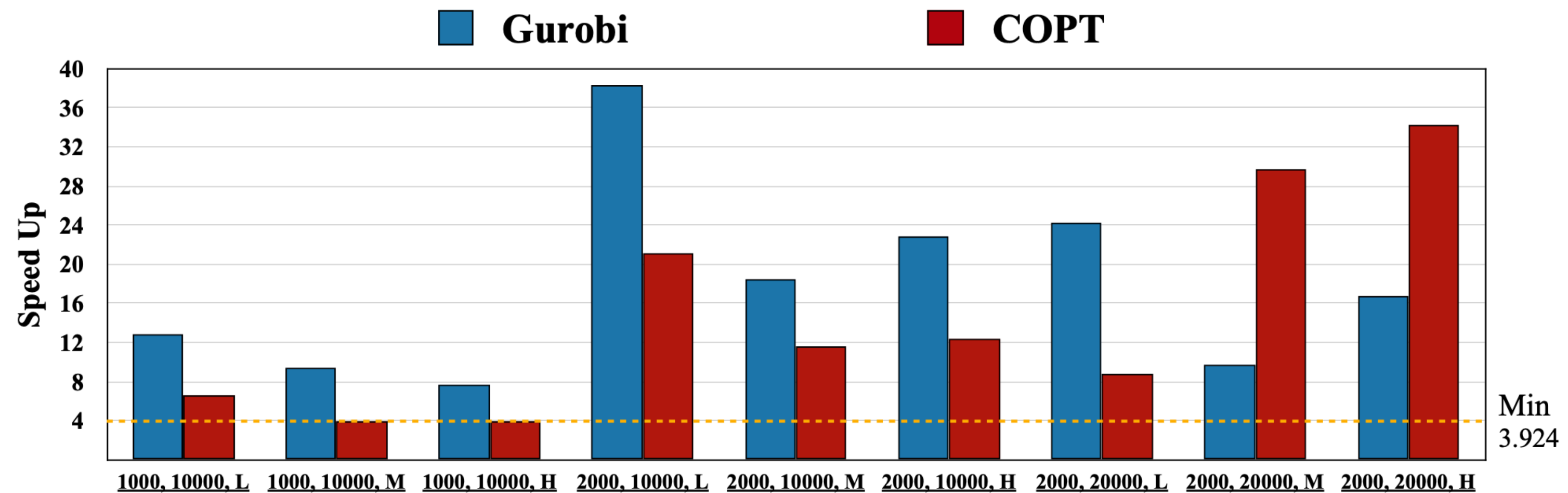
Average speedup on knapsack instances

# Numerical Test Results II

- Acceleration by two lines of code
- Remarkable speedup on primal solution finding for both the state of art MIP solvers Gurobi and COPT
- No loss of optimality



Unit Commitment



Set-Covering

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# Accelerated Second-Order Methods for Unconstrained Optimization and Applications

$$\min f(x), x \in X \text{ in } \mathbb{R}^n,$$

- where  $f$  is nonconvex and twice-differentiable,

$$g_k = \nabla f(x_k), H_k = \nabla^2 f(x_k)$$

- Goal: find  $x_k$  such that:

$$\|g_k\| \leq \epsilon \quad (\text{primary, first-order condition})$$

$$\lambda_{\min}(H_k) \geq -\sqrt{\epsilon} \quad (\text{secondary, second-order condition})$$

- First-order methods (FOM) typically need  $\mathbf{O}(\mathbf{n}^2\epsilon^{-2})$  arithmetic operations
- Second-order methods (SOM) typically need  $\mathbf{O}(\mathbf{n}^3\epsilon^{-1.5})$  arithmetic operations
- New? Yes, HSODM and DRSOM: a single-loop method with  $\mathbf{O}(\mathbf{n}^2\epsilon^{-1.75})$  operations  
(Zhang et al. <https://arxiv.org/abs/2211.08212>)

# An Integrated Descent Direction Using the Homogeneous Model (Zhang et al. SHUFE, 2022)

- The Newton-based trust-region method minimizes the Taylor quadratic model

$$\begin{aligned} \min_{d \in \mathbb{R}^n} m_k(d) &:= g_k^T d + \frac{1}{2} d^T H_k d \\ \text{s.t. } \|d\| &\leq \Delta_k. \end{aligned} \quad \longrightarrow \quad \psi_k(\xi_0, t; \delta) := \frac{1}{2} \begin{bmatrix} \xi_0 \\ t \end{bmatrix}^T \begin{bmatrix} H_k & g_k \\ g_k^T & -\delta \end{bmatrix} \begin{bmatrix} \xi_0 \\ t \end{bmatrix} = \frac{t^2}{2} \begin{bmatrix} \xi_0/t \\ 1 \end{bmatrix}^T \begin{bmatrix} H_k & g_k \\ g_k^T & -\delta \end{bmatrix} \begin{bmatrix} \xi_0/t \\ 1 \end{bmatrix}$$

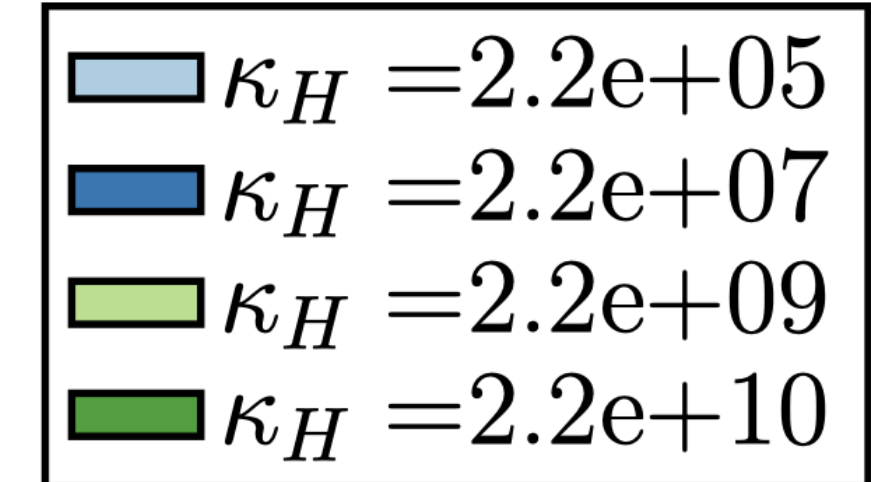
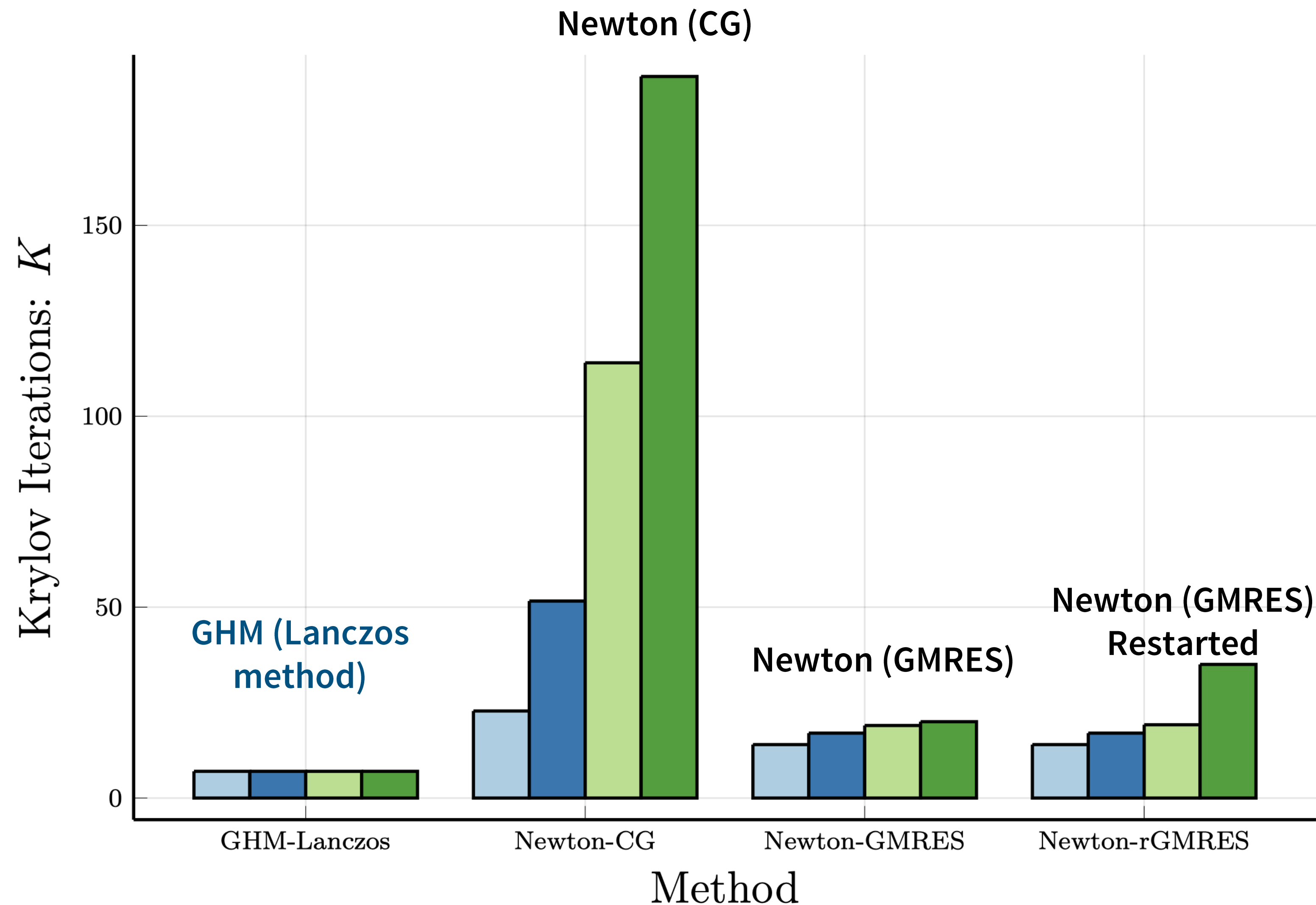
where  $\Delta_k = \epsilon^{1/2} / M$  is the trust radius

- $-g_k$  is the first-order steepest descent direction but ignores Hessian; the most-left eigenvector of  $H_k$ -would be a descent direction for the second order term
- Could we construct a direction integrating both?

**Answer:** with a suitable  $\delta_k$  and use the most-left eigen vector as the direction to go – a single loop algorithm to solve the original problem and replace the Newton step.

**$O(n^2 \epsilon^{-1.75})$  vs  $O(n^3)$  operations**

# Advantage of the Homogeneous Direction over Newton's



- Consider Hilbert matrix
- Computing a **Newton step** by iterative methods versus **Homogeneous Direction** from  $H + \delta I$
- $\delta \downarrow$  Condition number  $\kappa_H \uparrow$

# Dimension Reduced Second-Order Method (DRSOM)

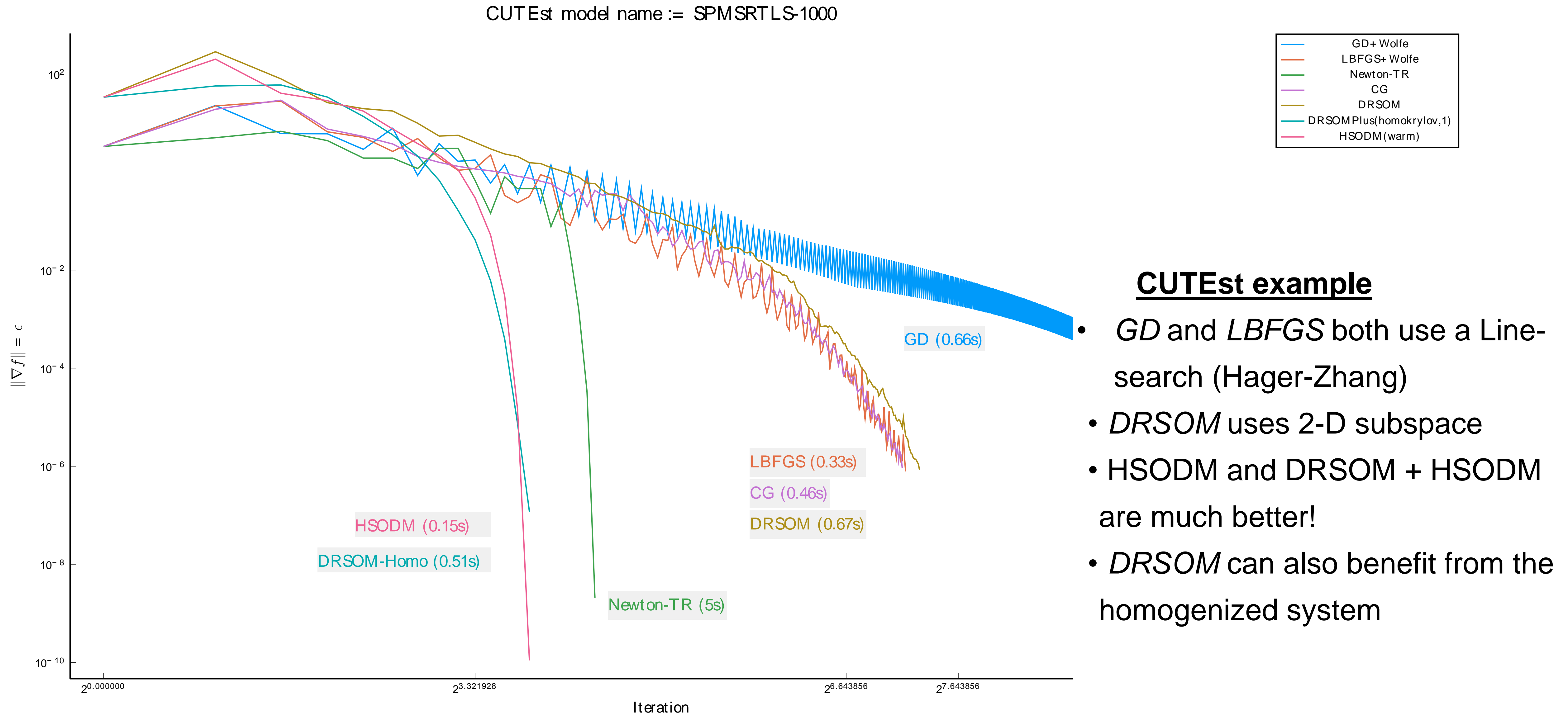
- Motivation from Multi-Directional FOM and Subspace Method, such as CG and ADAM, DRSOM applies the trust-region method in low dimensional subspace.
  - This results in a low-dimensional quadratic sub-minimization problem:
  - Typically, DRSOM adopts two directions  $d = -\alpha^1 \nabla f(x_k) + \alpha^2 d_k$
- where  $g_k = \nabla f(x_k)$ ,  $H_k = \nabla^2 f(x^k)$ ,  $d_k = x_k - x_{k-1}$
- Then we solve a 2-d quadratic minimization problem to decide the two step-sizes:

$$\min m_k^\alpha(\alpha) := f(x_k) + (c_k)^T \alpha + \frac{1}{2} \alpha^T Q_k \alpha$$

$$G_k = \begin{bmatrix} \|\alpha\|_{G_k} \leq \Delta_k \\ g_k^T g_k & -g_k^T d_k \\ -g_k^T d_k & d_k^T d_k \end{bmatrix}, Q_k = \begin{bmatrix} g_k^T H_k g_k & -g_k^T H_k d_k \\ -g_k^T H_k d_k & d_k^T H_k d_k \end{bmatrix}, c_k = \begin{bmatrix} -\|g_k\|^2 \\ g_k^T d_k \end{bmatrix}$$



# Preliminary Results: HSODM, DRSOM and DRSOM+HSODM



# Sensor Network Localization I

- Consider Sensor Network Location (SNL)

$$N_x = \{(i, j) : \|x_i - x_j\| = d_{ij} \leq r_d\}, N_a = \{(i, k) : \|x_i - a_k\| = d_{ik} \leq r_d\}$$

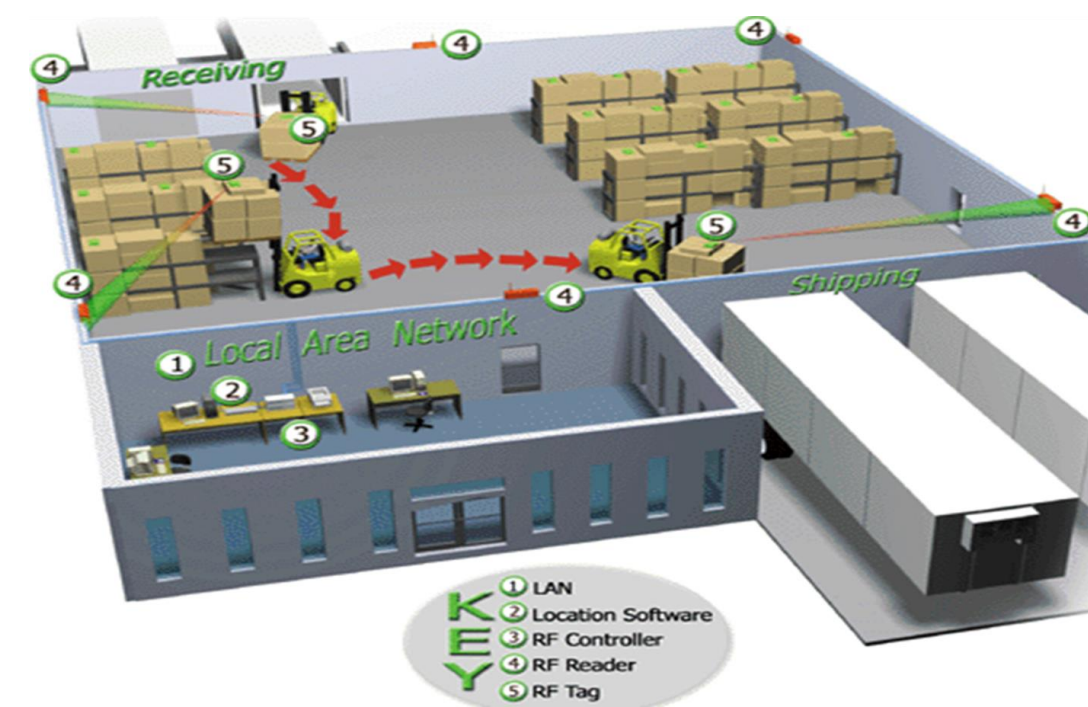
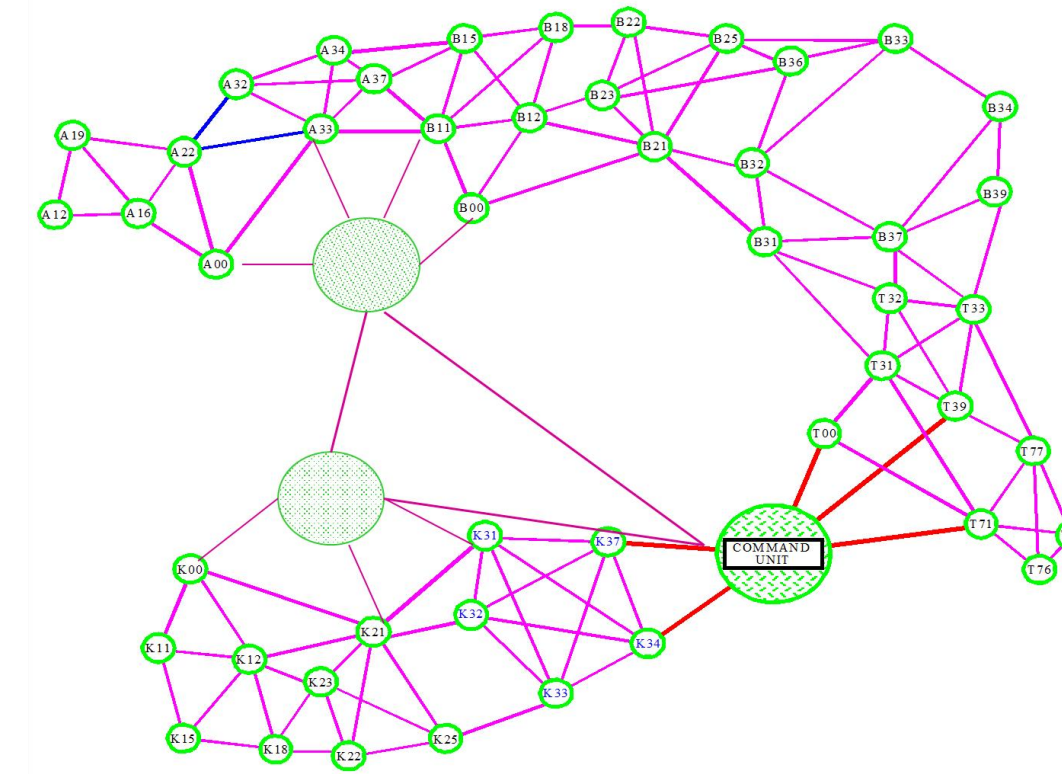
where  $r_d$  is a fixed parameter known as the radio range. The SNL problem considers the following QCQP feasibility problem,

$$\|x_i - x_j\|^2 = d_{ij}^2, \forall (i, j) \in N_x$$

$$\|x_i - a_k\|^2 = \bar{d}_{ik}^2, \forall (i, k) \in N_a$$

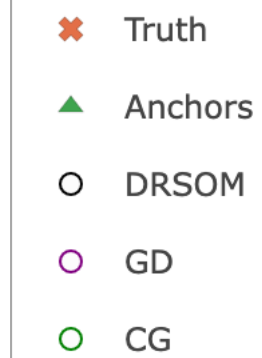
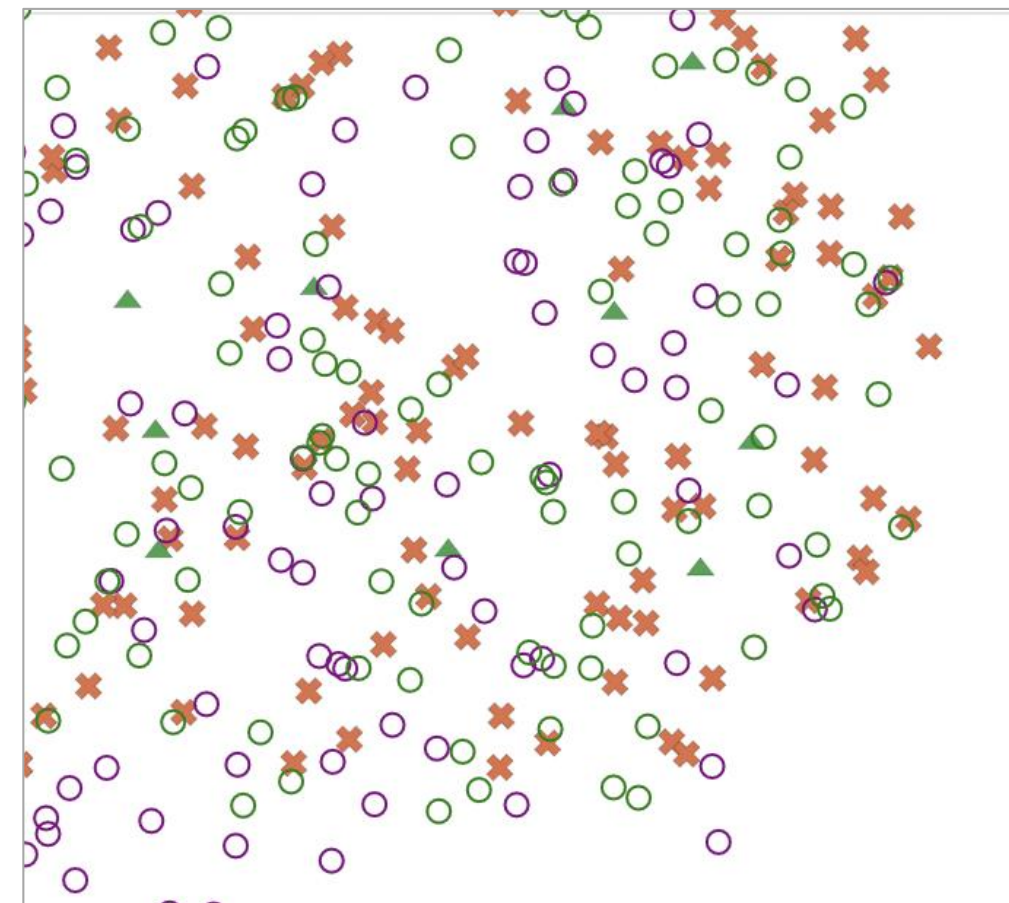
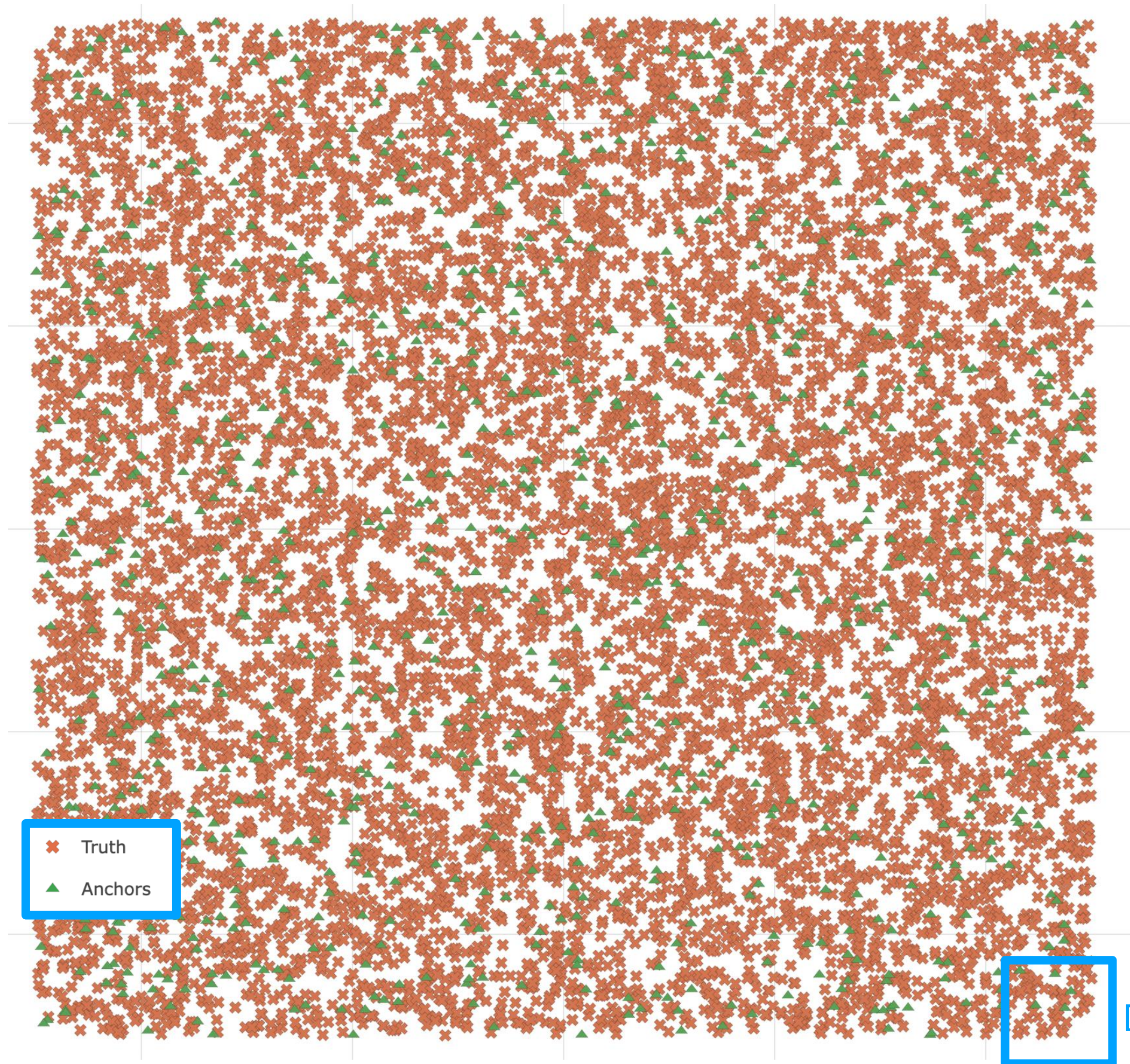
- We can solve SNL by the nonconvex nonlinear least square (NLS) problem

$$\min_X \sum_{(i,j) \in N_x} (\|x_i - x_j\|^2 - d_{ij}^2)^2 + \sum_{(k,j) \in N_a} (\|a_k - x_j\|^2 - \bar{d}_{kj}^2)^2.$$

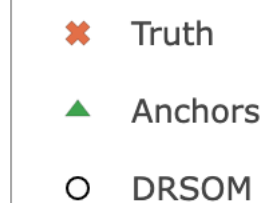
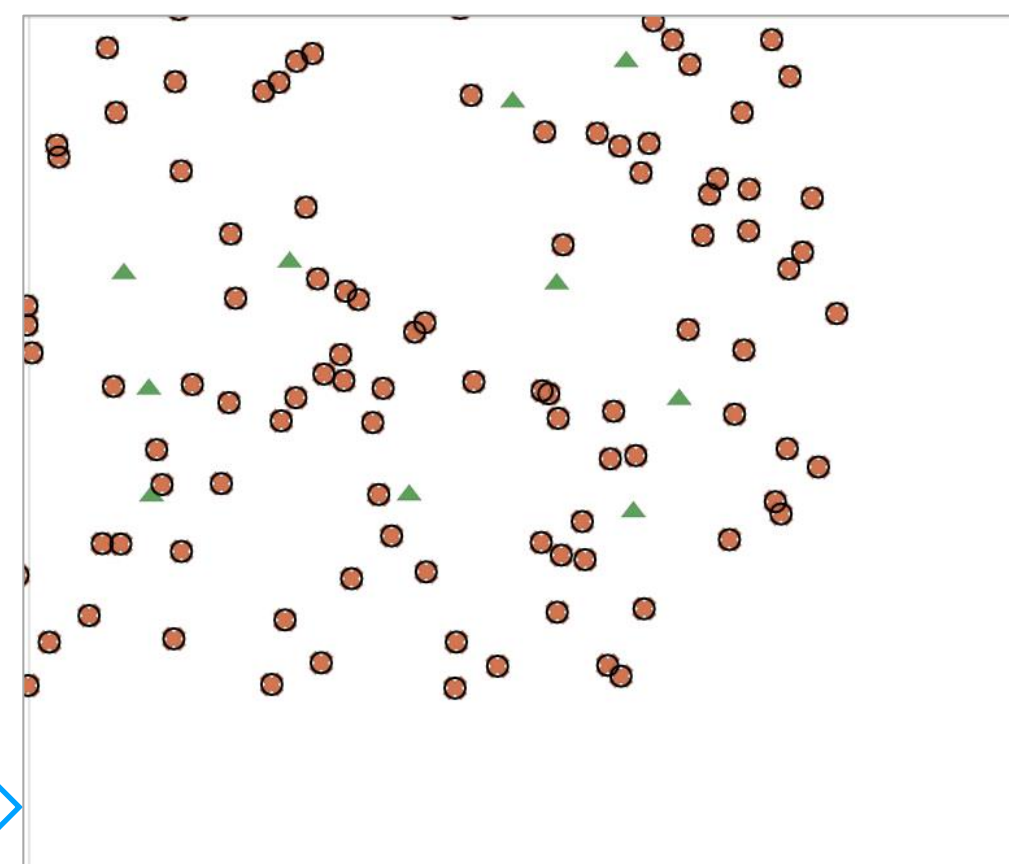


# Sensor Network Location II

- Graphical results with 10,000 nodes and 1000 anchors (no noise) **within 3,000 seconds**



- GD with Line-search and Hager-Zhang CG both timeout**



- DRSOM can converge to  $|g_k| \leq 1e^{-5}$  in 2,200s**

# Zero-Order Optimization: SOLNP+

- First proposed by Y in 1989.
- Originally implemented (SOLNP) in **Matlab**, 1989.
- **R** implementation (**Rsolnp**) by Alexios Ghalanos and Stefan Theussl, 2011.
- **New and C implementation (SOLNP+) with improvements, 2022; and addition of Randomized BCG and DRSSOM for unconstrained optimization by Tan et al., 2023**
- **Github link: [https://github.com/COPT-Public/SOLNP\\_plus](https://github.com/COPT-Public/SOLNP_plus)**
- Use **forward difference** to evaluate the gradient.

$$[\nabla_{\delta} f(x)]_i = \frac{f(x + \delta e_i) - f(x)}{\delta}, \quad e_i = [0, \dots, 1, \dots, 0].$$

# RMP: Multi-Point Random Perturbation

- Multipoint ZO Gradient Estimates

$$\hat{\nabla} f(x) := \frac{\phi(n)}{\delta b} \sum_{i=1}^b [(f(x + \delta u_i) - f(x)) u_i]$$

$$\mathbb{E}[\hat{\nabla} f(x)] = \nabla f_{\delta}(x)$$

$$f_{\delta}(x) = \mathbb{E}_u[f(x + \delta u)]$$

where  $u_i$  is i.i.d. **random direction**.

- Advantage: **Fewer** function queries to evaluate the gradient.

# SOLNP+: Adopt Two Strategies

- With gradient estimates, SOLNP+ implements ZO version of
  - ZO-RMP (Ghadimiet al., 2013; Duchi et al., 2014), or
  - ZO-BCD, (recent research see Sun/Y 2020, Cai et al., 2021), that is, use  $e_i$  (vector with zero components except that dimension  $i$  is 1),  $i$  is randomly chosen
- DRSOM (Zhang et al., 2022) with interpolation

S. Ghadimi and G. Lan, “Stochastic first-and zeroth-order methods for nonconvex stochastic programming,” *SIAM J. Optimiz.*, vol. 23, no. 4, pp. 2341–2368, 2013. doi: 10.1137/120880811

J.C.Duchi, M. I. Jordan, M. J. Wainwright, and A. Wibisono, “Optimal rates for zero-order convex optimization: The power of two function evaluations,” *IEEE Trans. Inf Theory*, vol.61,no.5,pp.2788-2806,2015.doi: 10.1109/TIT.2015.2409256.

R. Sun and Y. Y, “Worst-case complexity of cyclic coordinate descent:  $O(n^2)$  gap with randomized version.” *Mathematical Programming*, Volume 185, 487-520, 2021.

Cai, HanQin, et al. "A zeroth-order block coordinate descent algorithm for huge-scale black-box optimization." *International Conference on Machine Learning*. PMLR, 2021.

Zhang, Chuwen, et al. "DRSOM: A Dimension Reduced Second-Order Method and Preliminary Analyses." *arXiv preprint arXiv:2208.00208* (2022)

# Experiments in Large Problems: Rosenbrock I

- Rosenbrock function is a well-known nonconvex functions in the form of

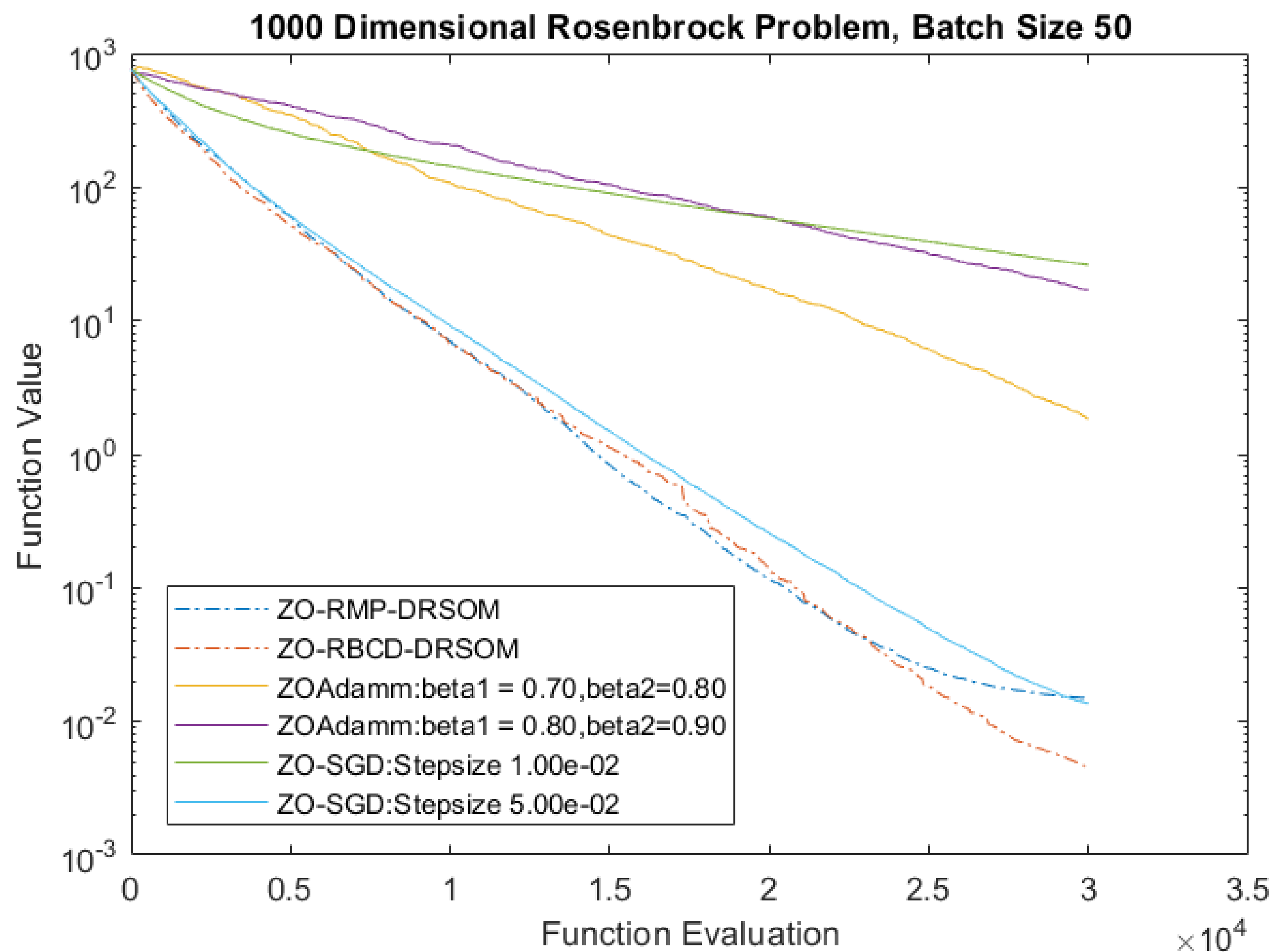
$$f(\mathbf{x}) = \sum_{i=1}^{n-1} 100(x_{i+1} - x_i^2)^2 + (x_i - 1)^2$$

- ZO-Adamm, ZO-SGD and ZO-DRSOM are tested in a 1200 dimensional Rosenbrock problem.
  - Batch size 50
  - Each experiment is repeated for 10 times.

Chen, Xiangyi, et al. "Zo-adamm: Zeroth-order adaptive momentum method for black-box optimization." *Advances in neural information processing systems* 32 (2019).

# Experiments in Large Problems: Rosenbrock II

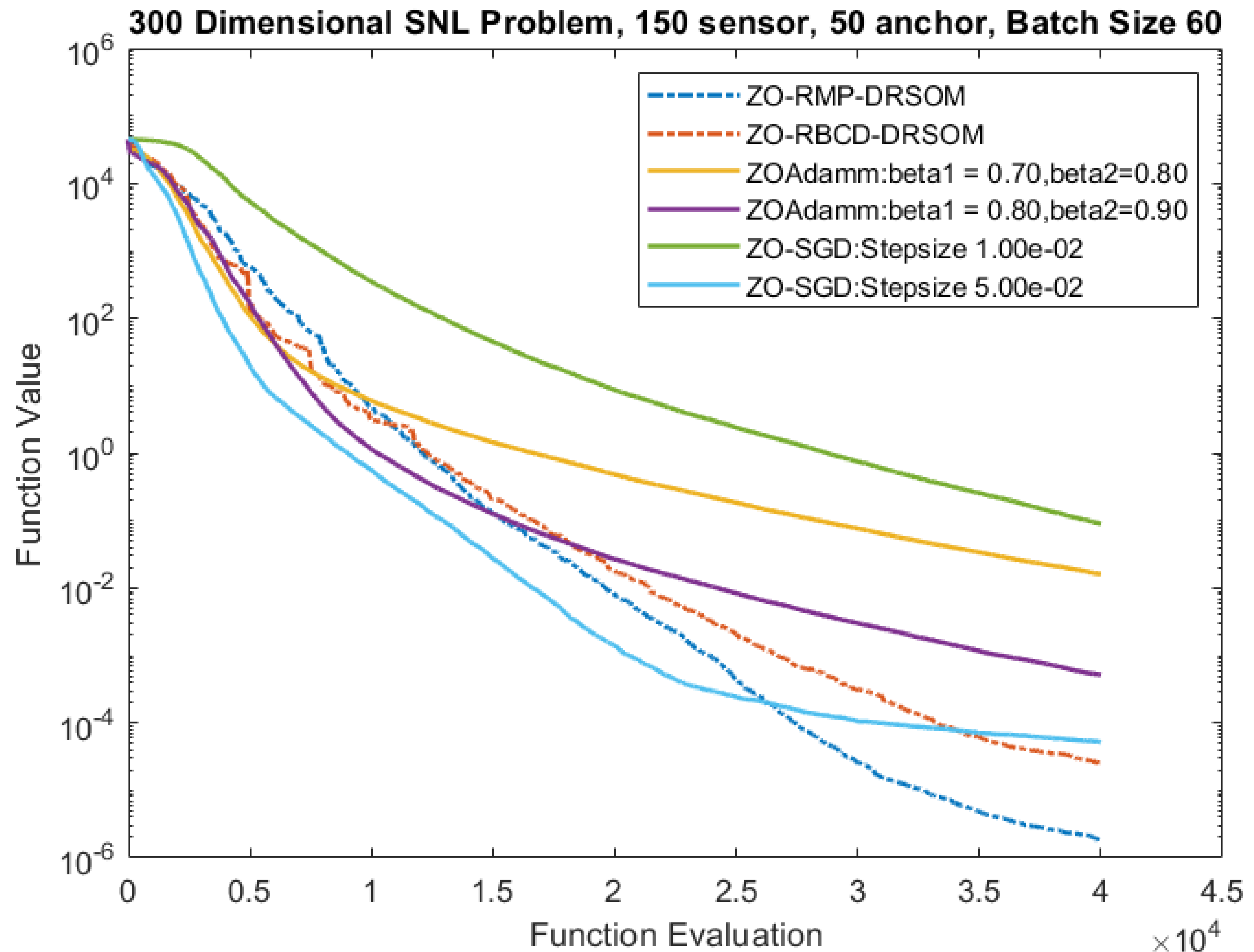
- ZO-RMP-DRSOM, ZO-RBCD-DRSOM and ZO-SGD decrease most smoothly. However, inappropriate parameters lead to worse performance of ZO-SGD and ZO-ADAMM.





# Experiments in Large-Scale Problems: SNL I

- ZO-Adam, ZO-SGD, ZO-RBCD-DRSOM and ZO-RMP-DRSOM are tested in a 150-sensor SNL problem.
  - Batch size 60
  - Each experiment is repeated for 10 times.



# Overall Takeaways

Know the pros and cons of OR and AI models and use them intelligently

Most OR optimization models/algorithms can be readily adopted for online training

Pre-training greatly improves Mixed Integer LP solvers that benefit real economy

Better to integrate ZOM, FOM and SOM for Nonlinear and/or Black-Box Optimization!

• **THANK YOU**